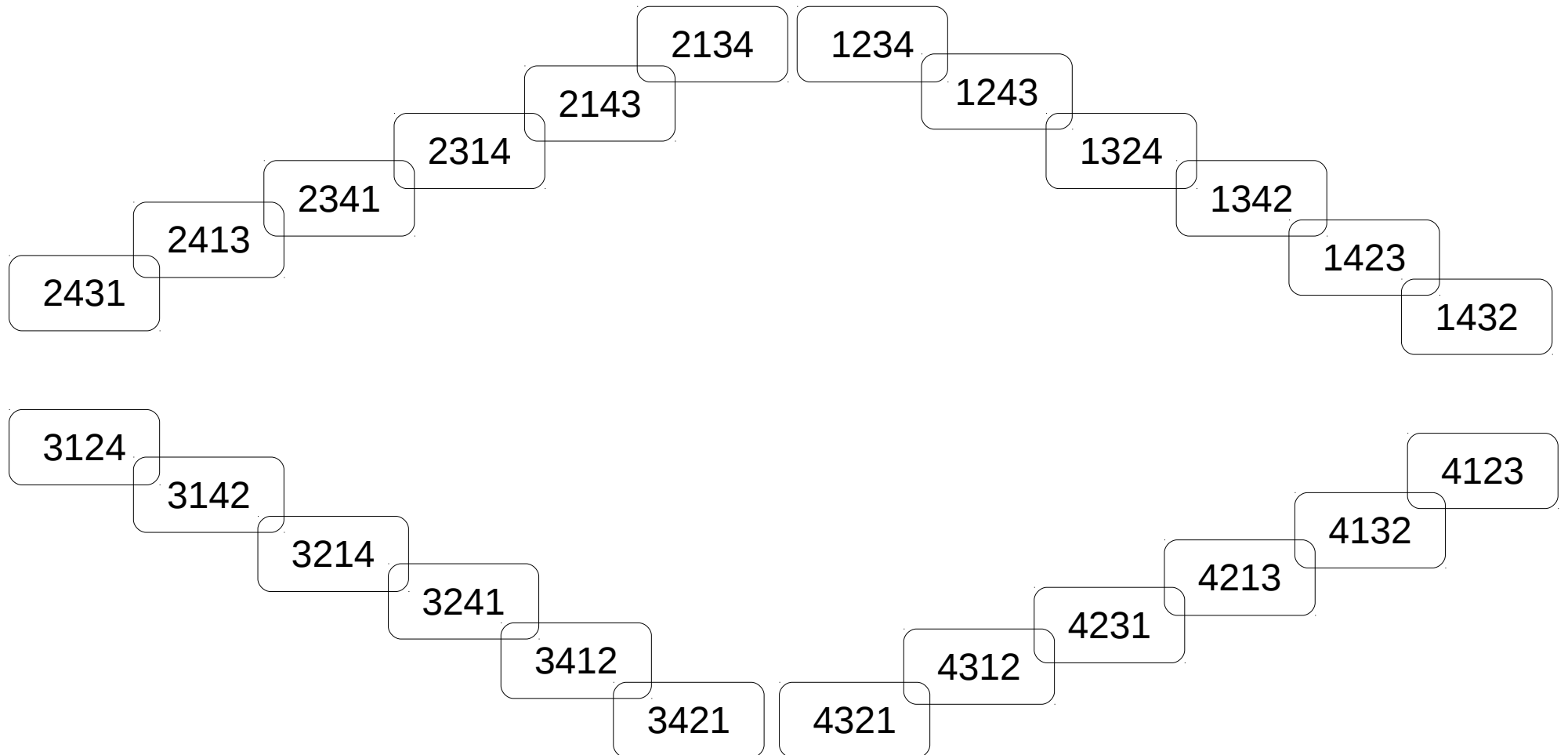


EARIN
Jarosław Arabas
Problem solving by searching

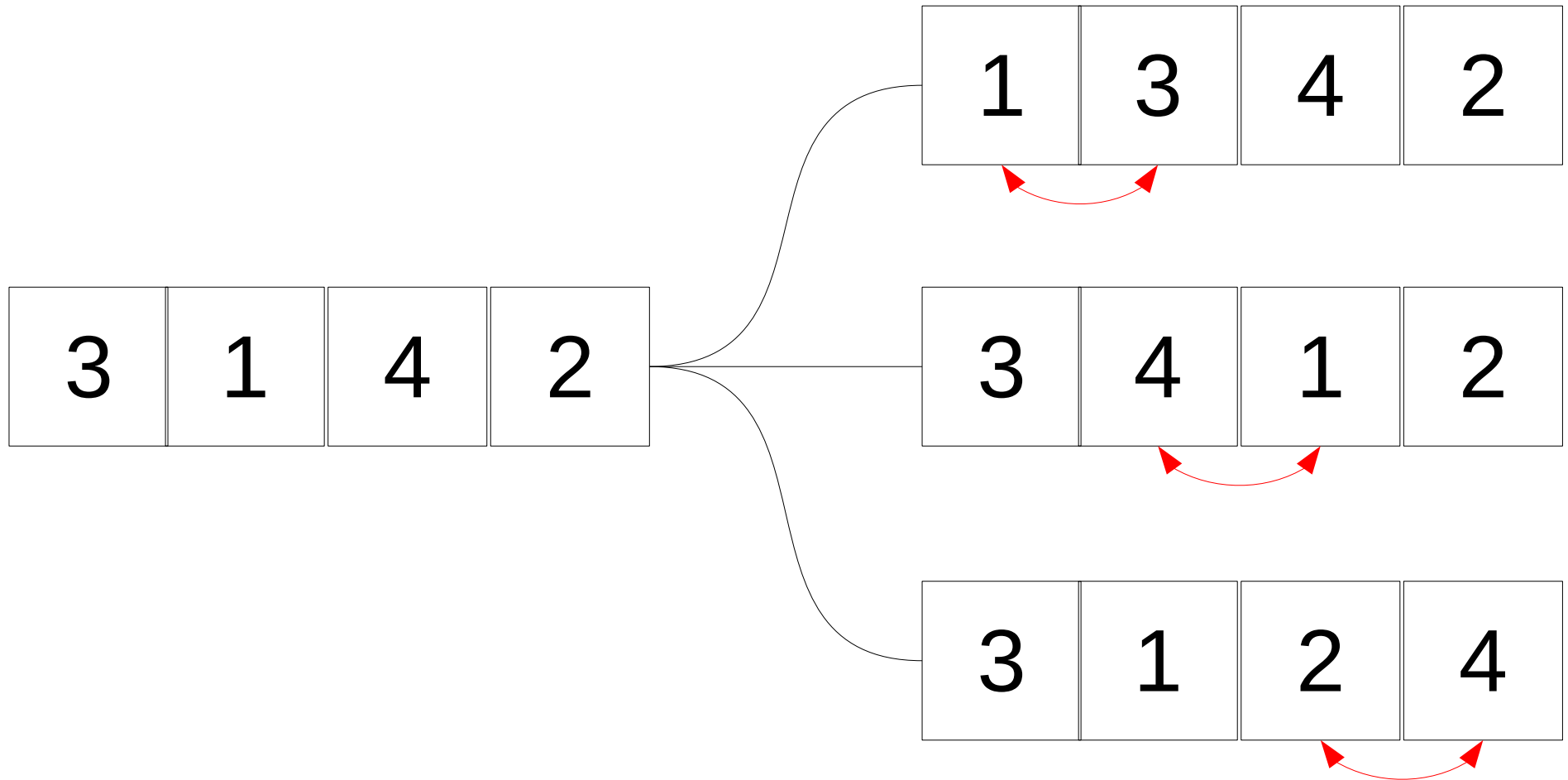
Sort a table of integers

```
void sort (int *tab, unsigned n){
    int i=1,j=0,sorted=0;
    do{
        sorted=1;
        for(j=0;j<n-i;++j){
            if (tab[j+1]<tab[j]){
                int temp=tab[j];
                tab[j]=tab[j+1];tab[j+1]=temp;
                sorted=0;}
            ++i;
        }while(i<n && !sorted);
    }
```

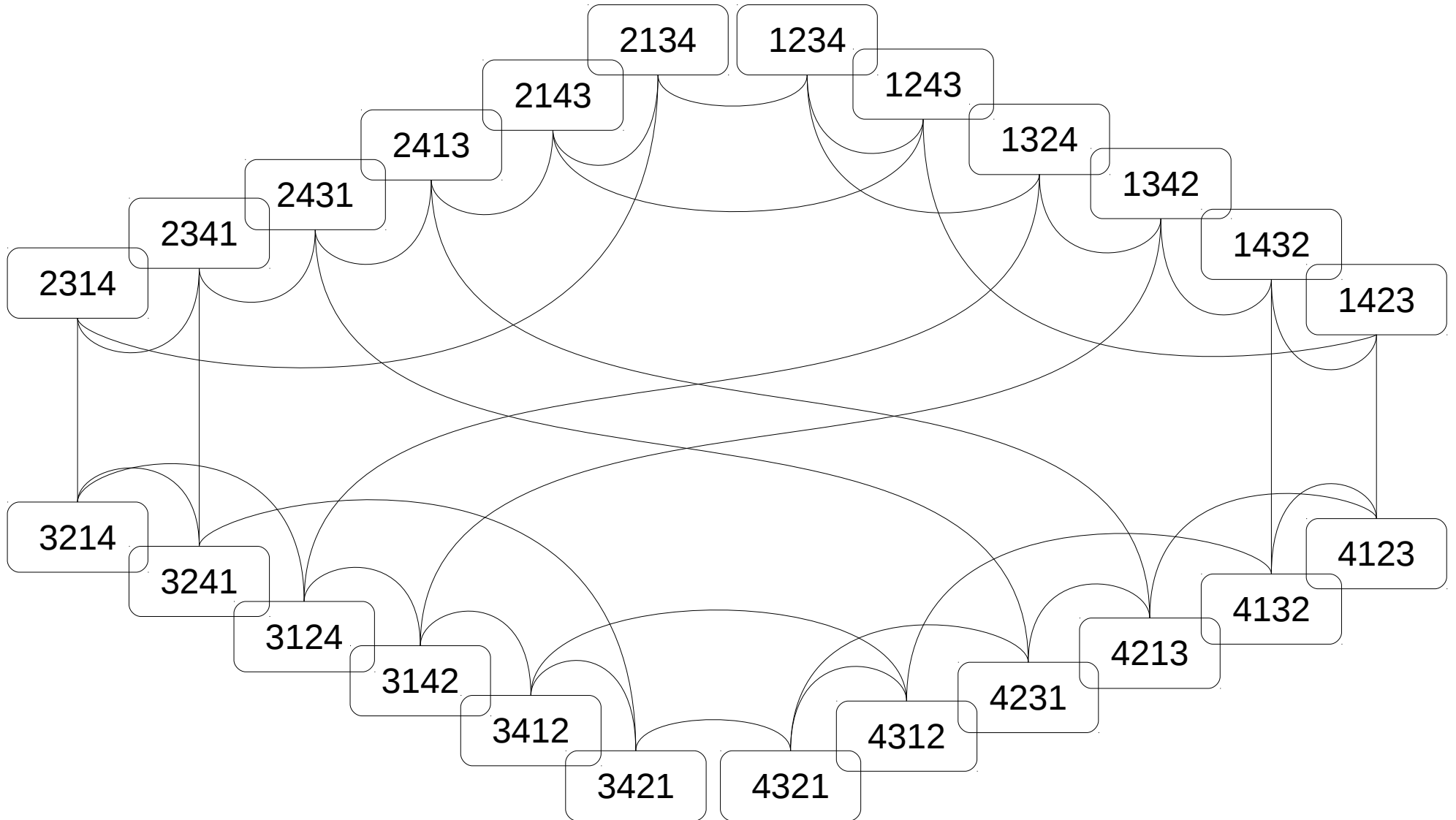
Set of possible table orderings



Relations between solutions

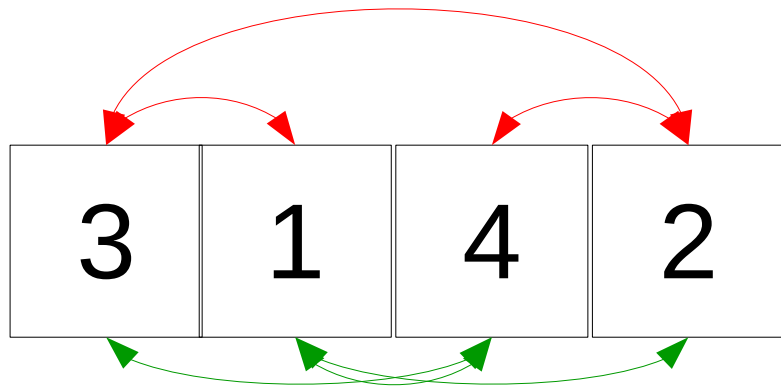


Search space



Objective function

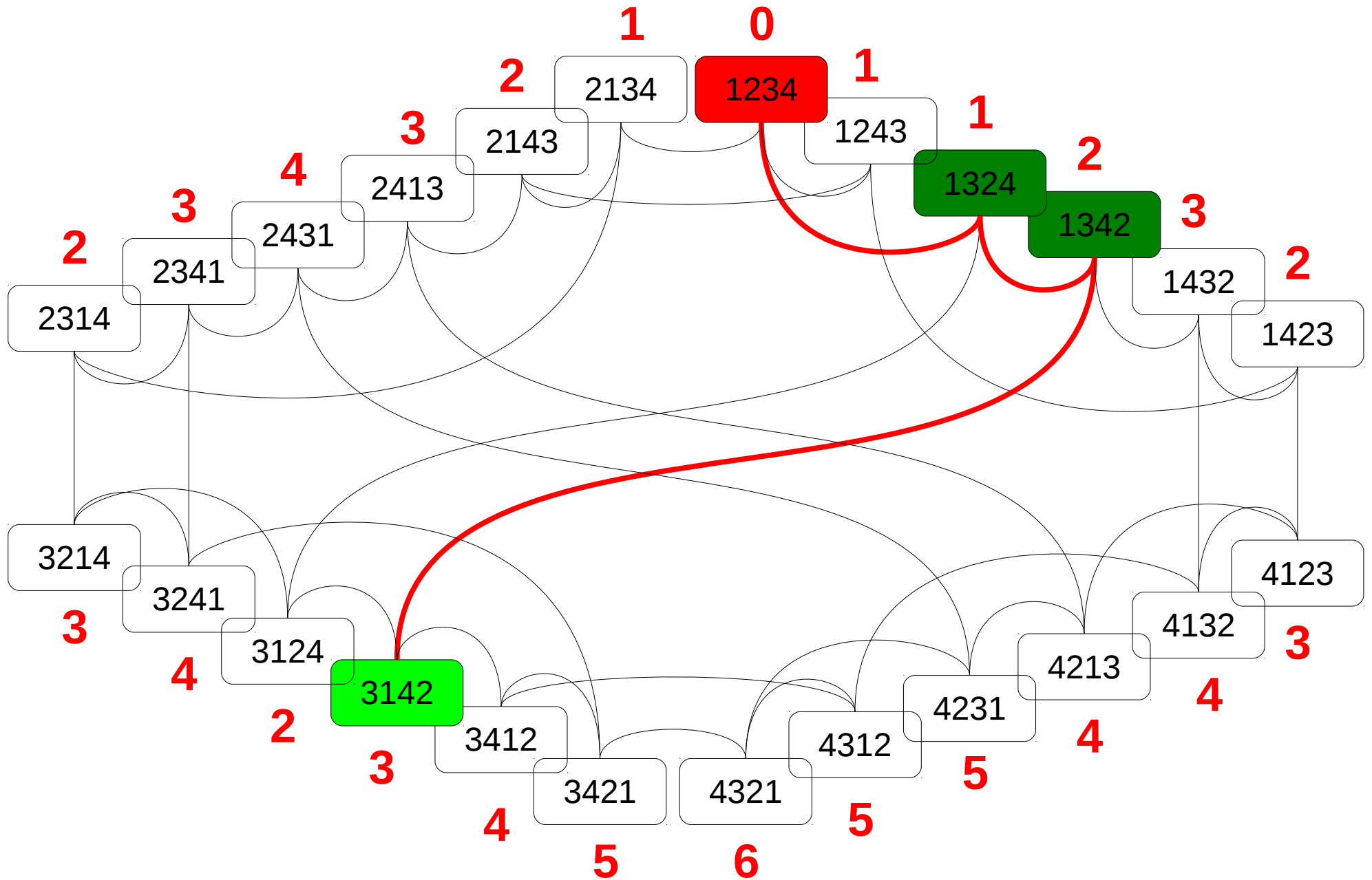
- Measure of the solutions quality



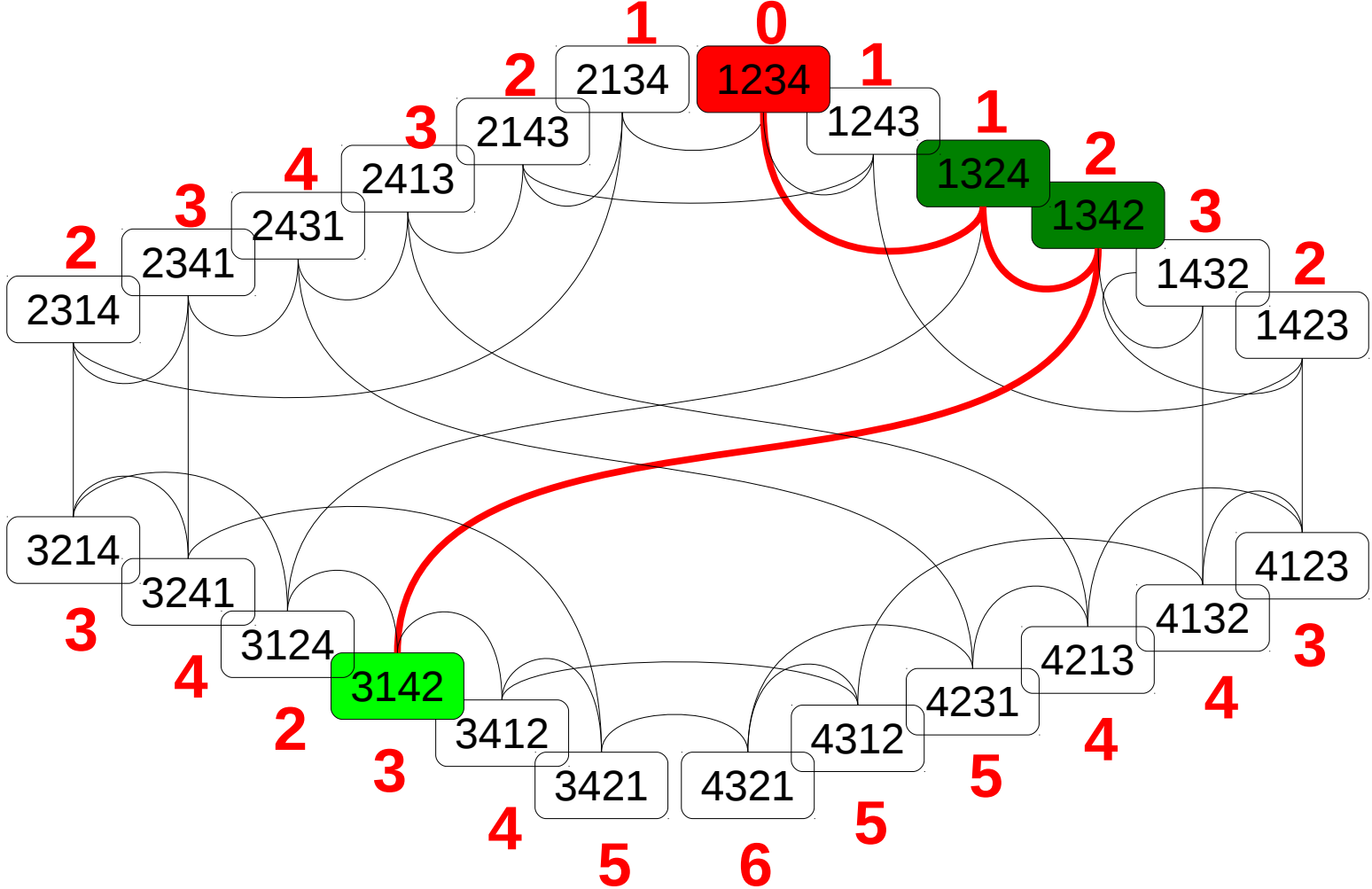
3 improperly ordered pairs

3 properly ordered pairs

Example path in the search space



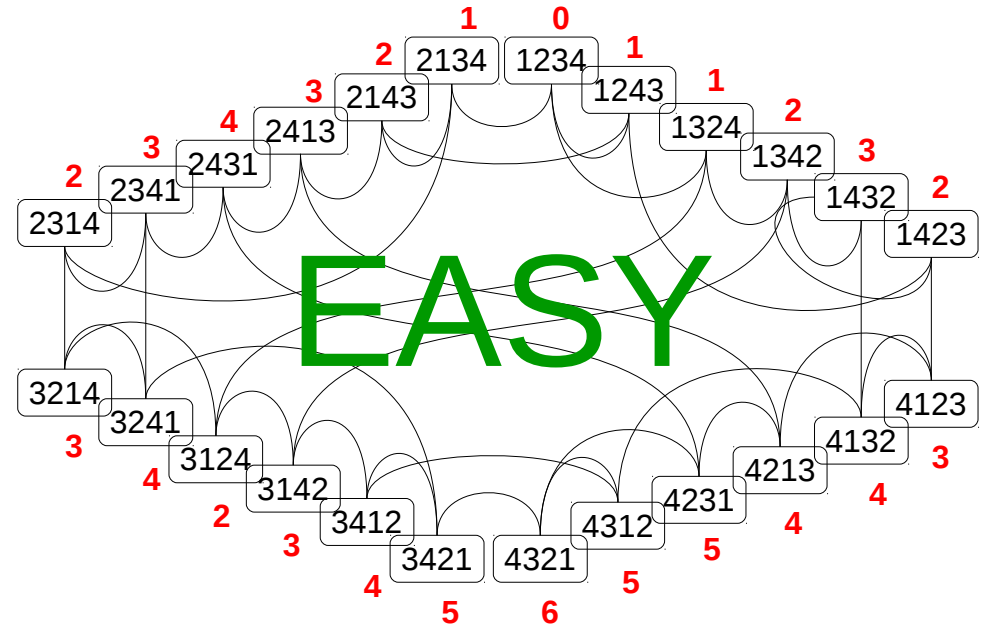
Path in the search space



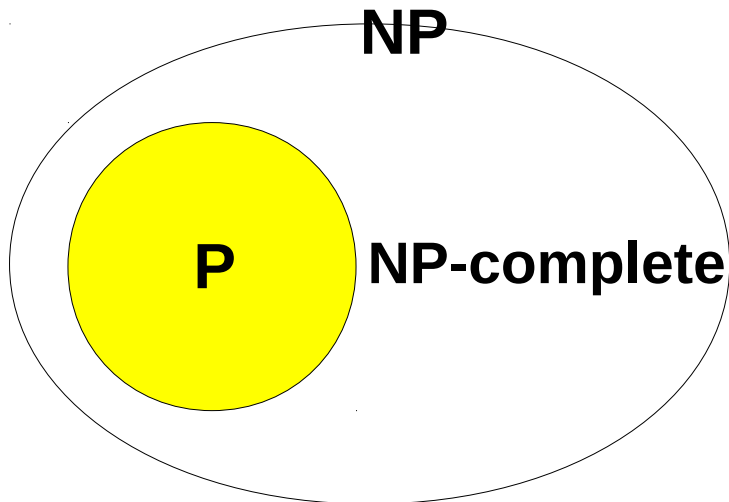
Search space
serialization

3142 → **1342** → **1324** → **1234**

“Easy” and “hard” problems

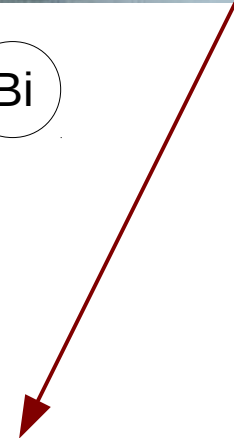
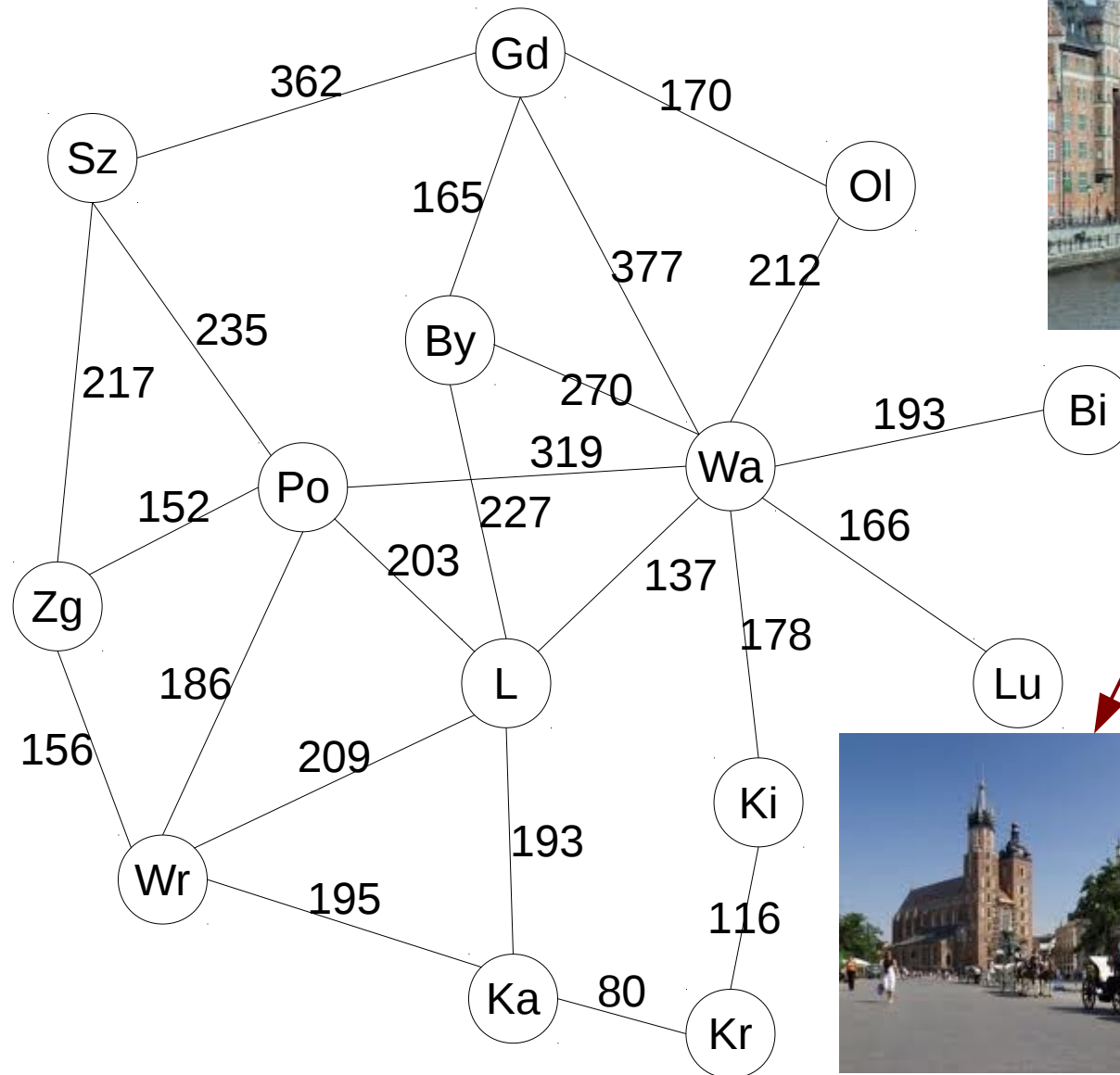


Hierarchy of decision problems

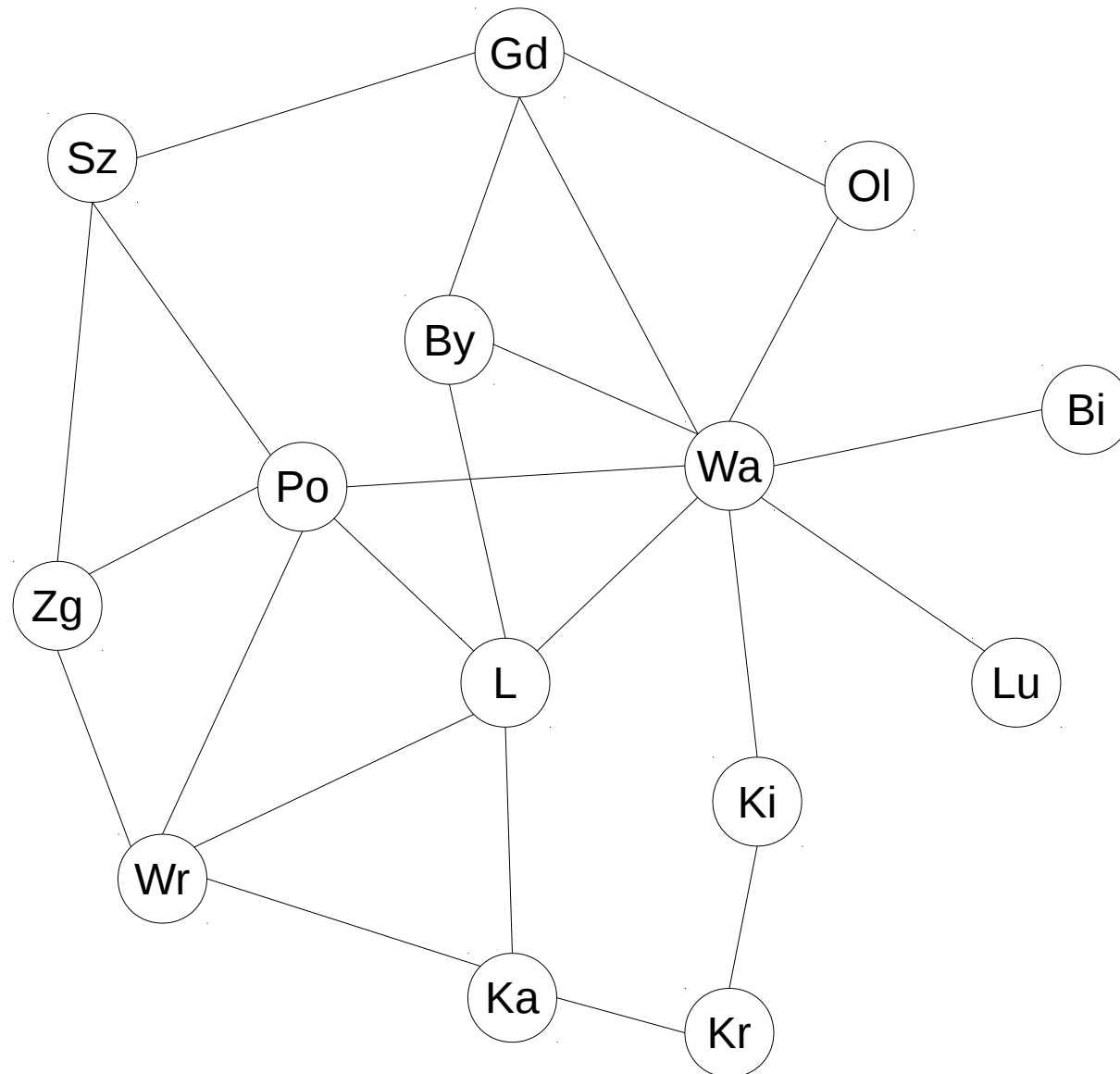


- Small neighborhoods (polynomial of n)
- Short paths (polynomial of n)
- Single optimum of the objective function

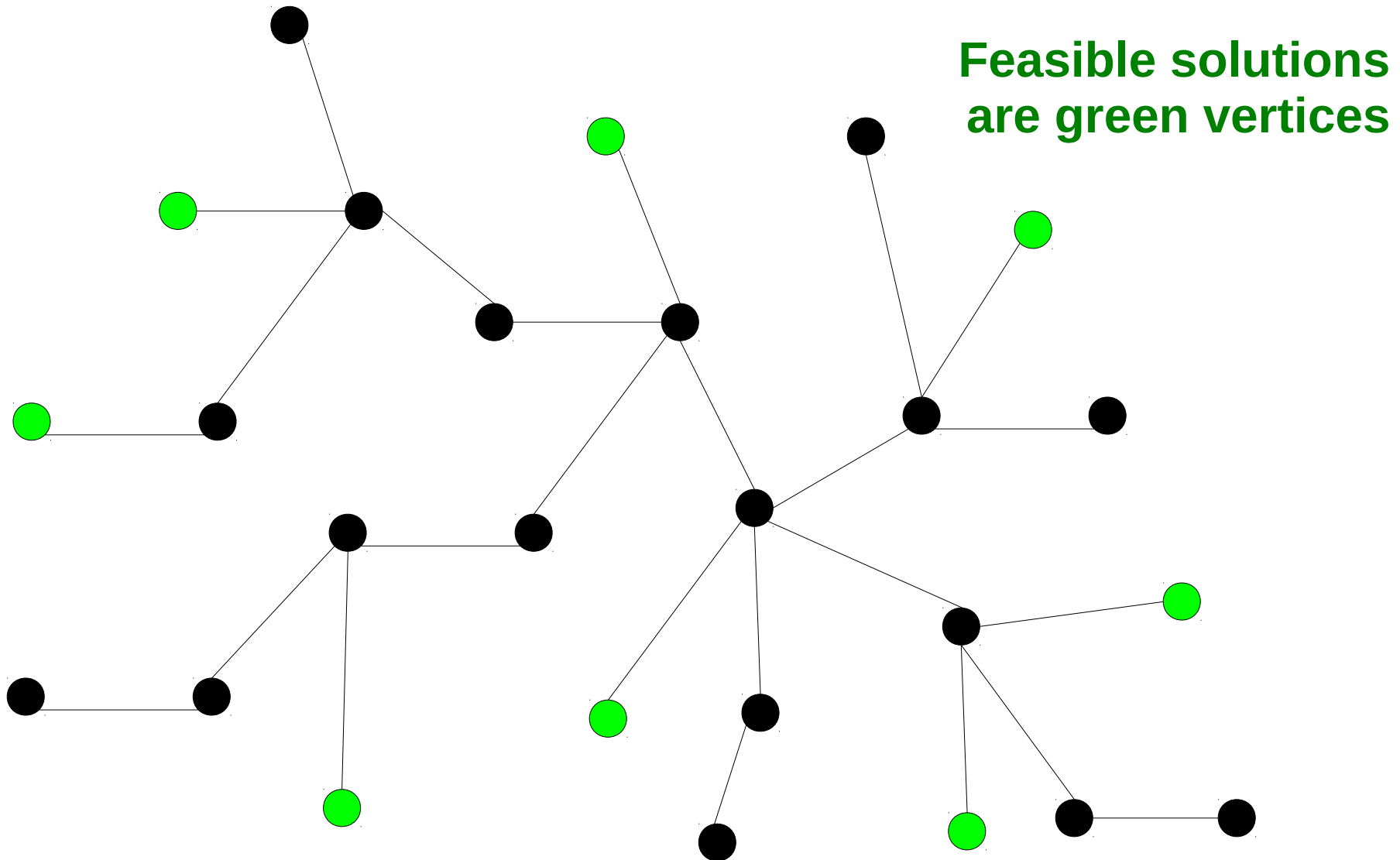
Example problem shortest trip from Gd to Kr



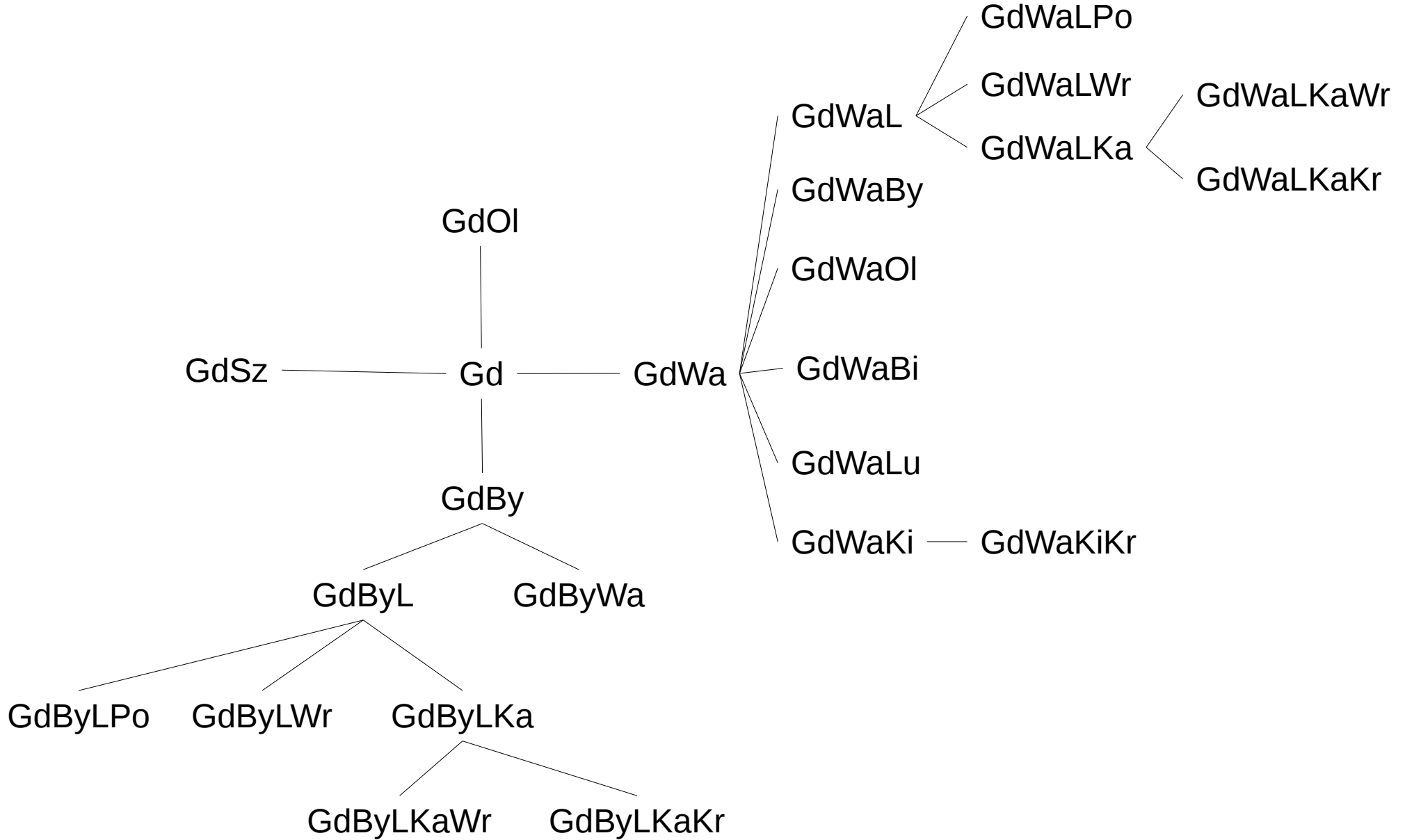
Uninformed search



Structure of the search space



Fragment of the search space



Breadth-first and depth-first search

procedure breadth first

$A \leftarrow \{s_0\}$

while $A \neq \emptyset$

$x \leftarrow \text{popFIFO}(A)$

$Y \leftarrow \text{neighbors}(x)$

$A \leftarrow A \cup Y$

procedure depth first

$A \leftarrow \{s_0\}$

while $A \neq \emptyset$

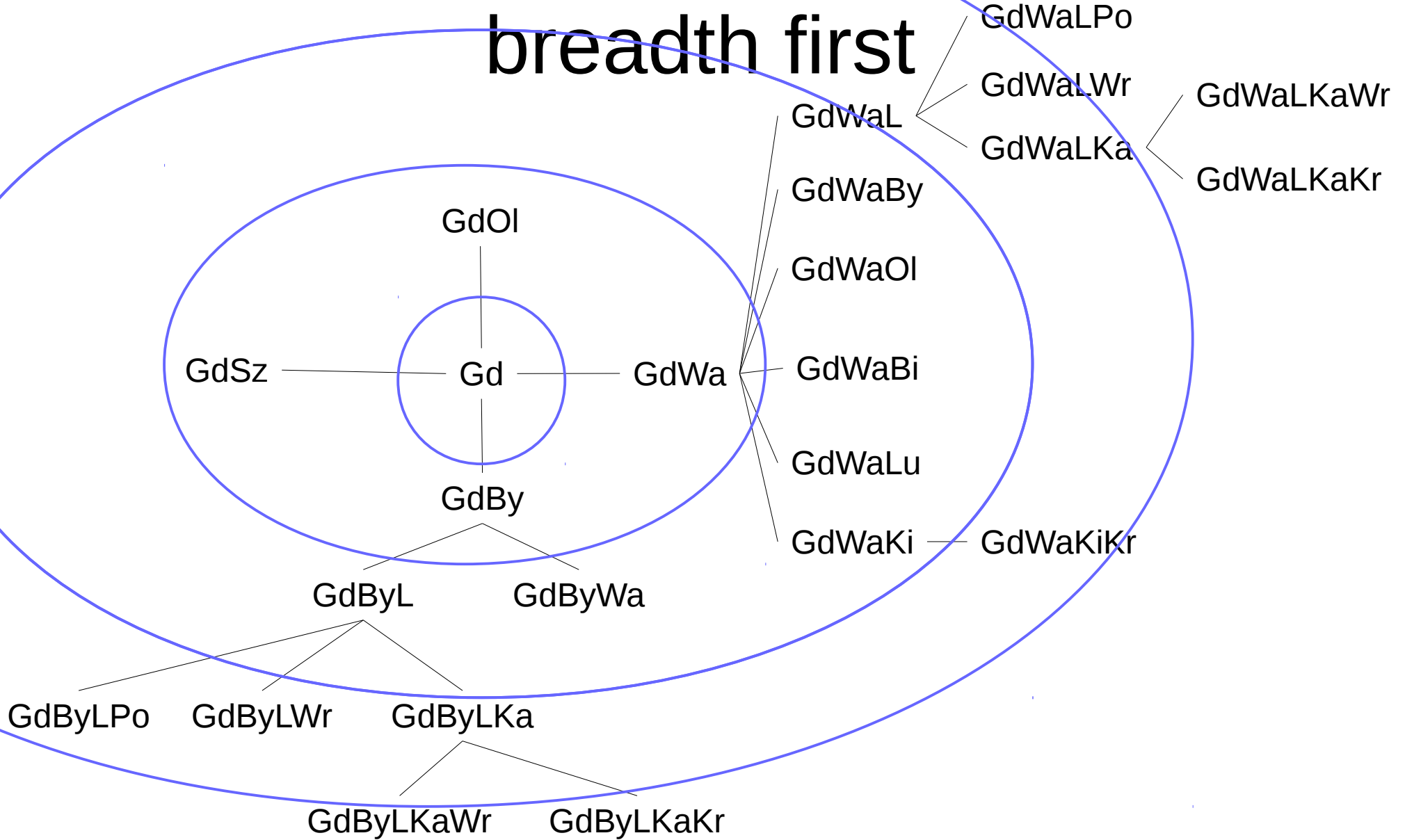
$x \leftarrow \text{popLIFO}(A)$

$Y \leftarrow \text{neighbors}(x)$

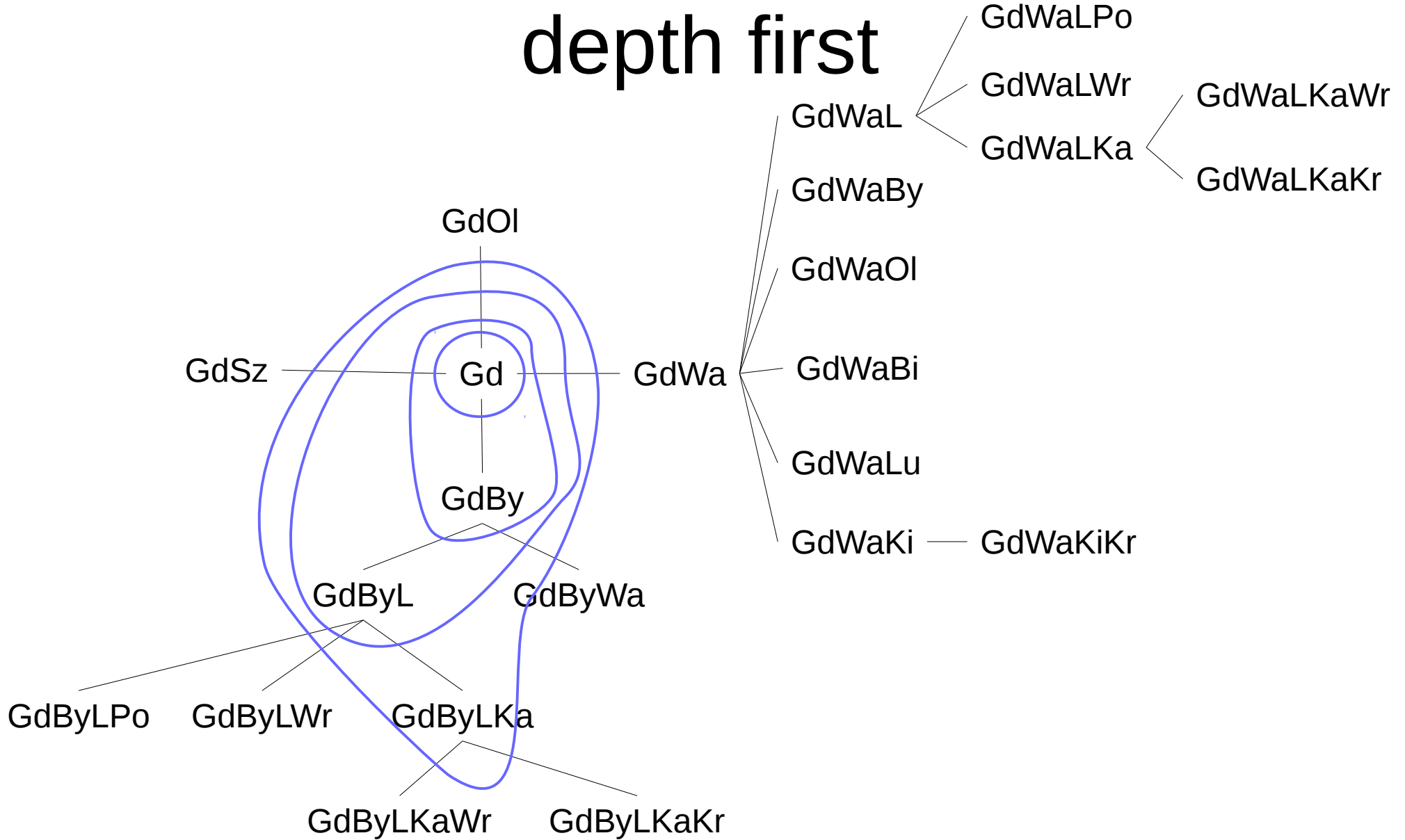
$A \leftarrow A \cup Y$

Sequence of visited nodes

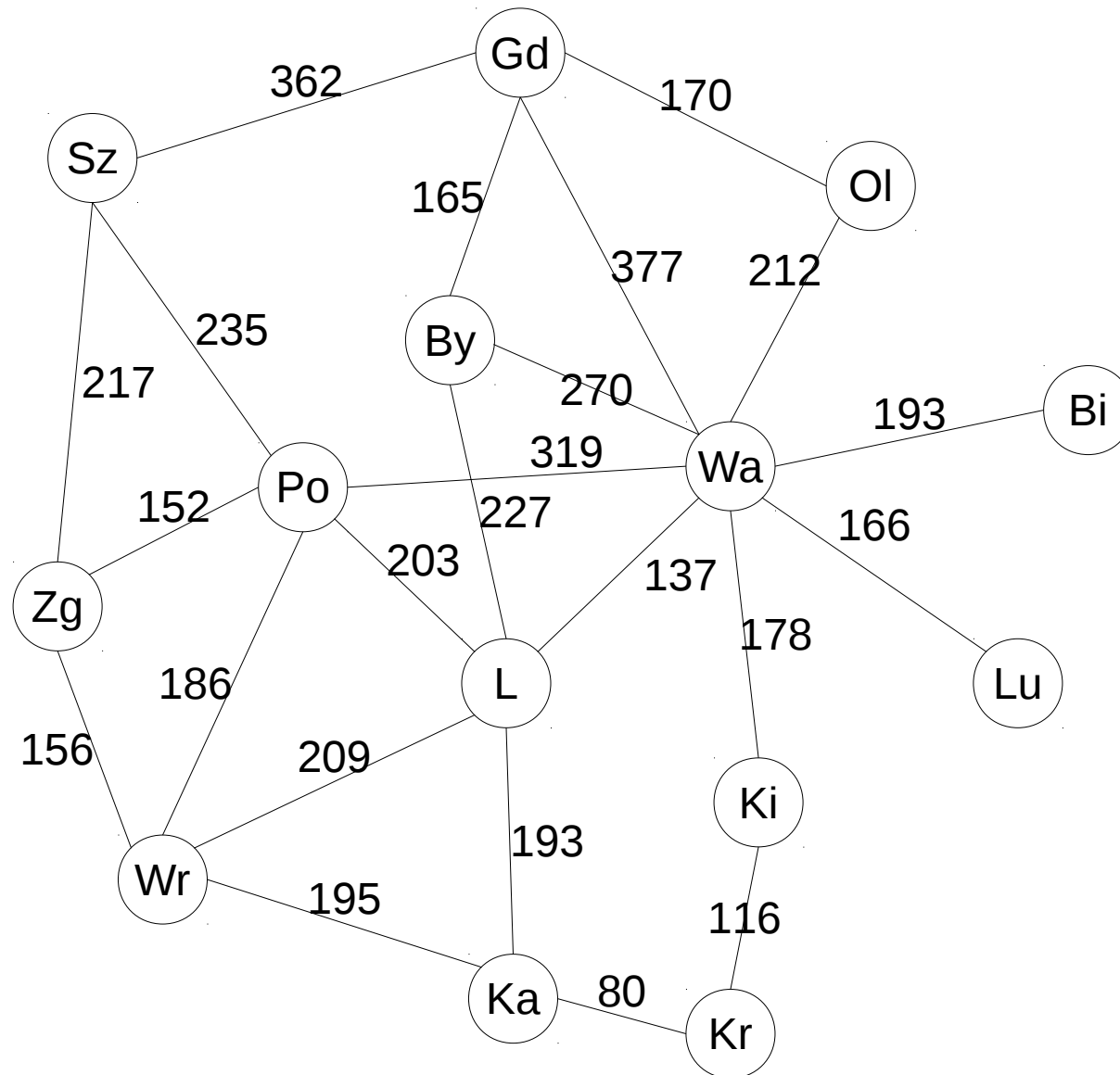
breadth first



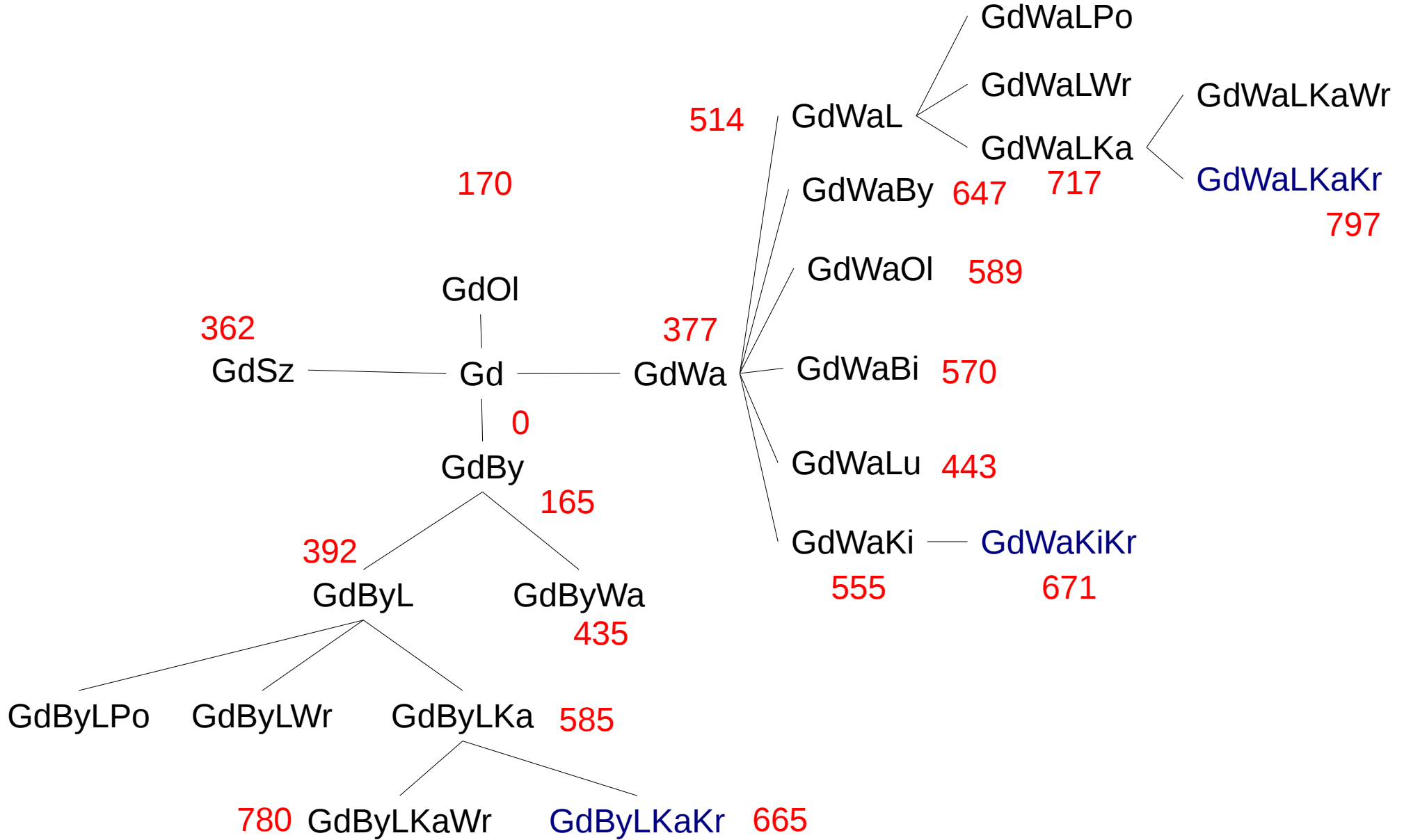
Sequence of visited nodes depth first



Informed search



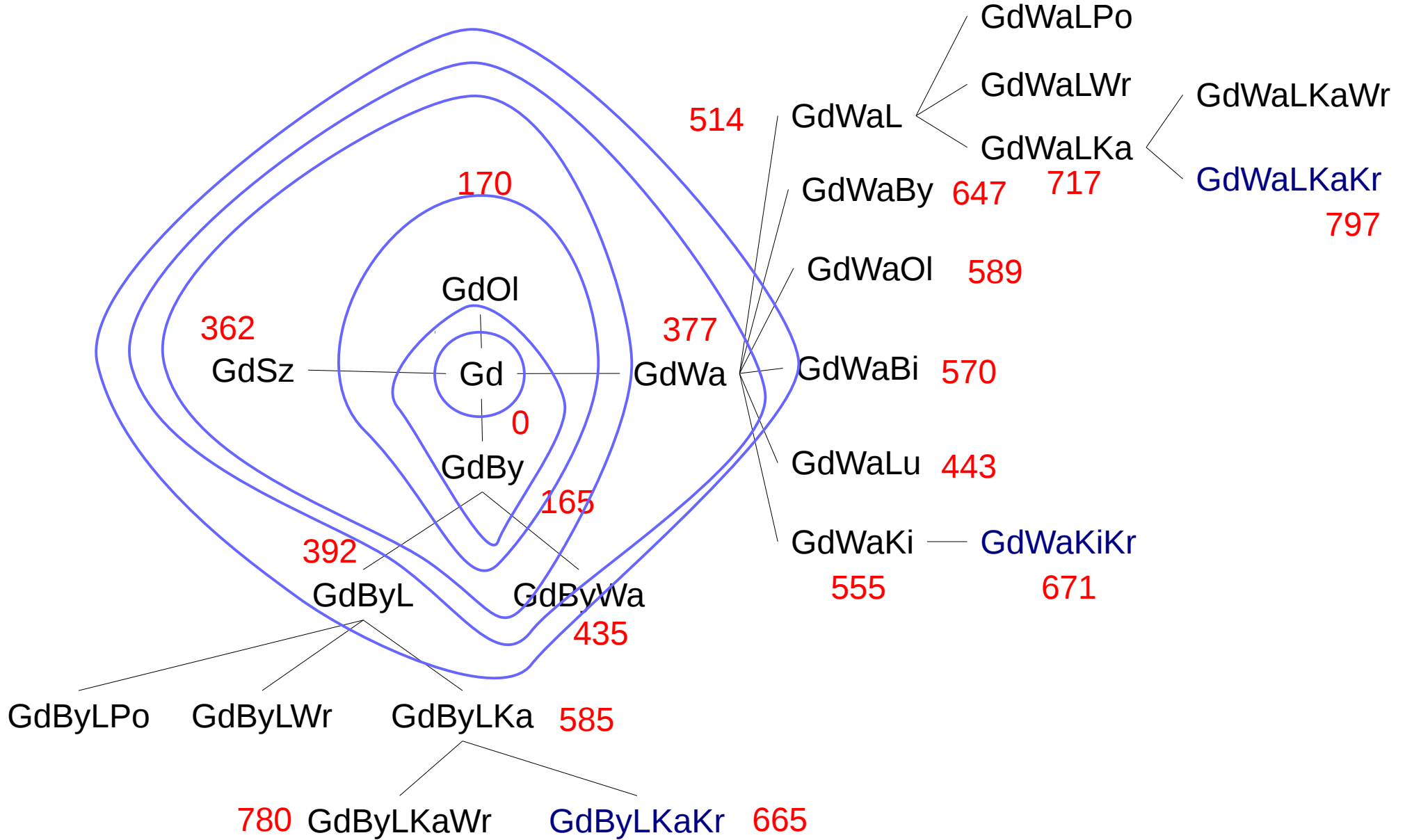
Fragment of the search space

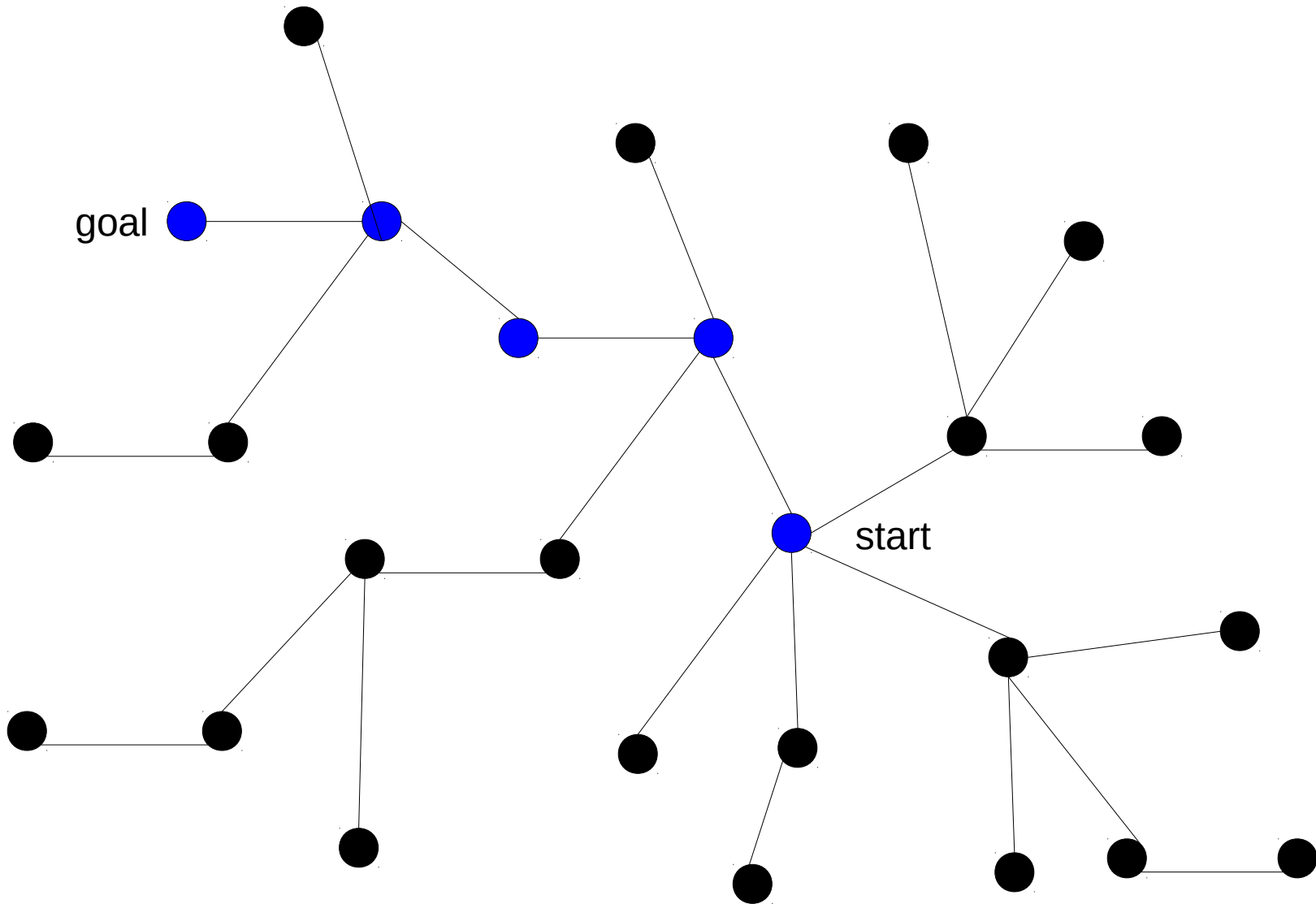


Best-first search

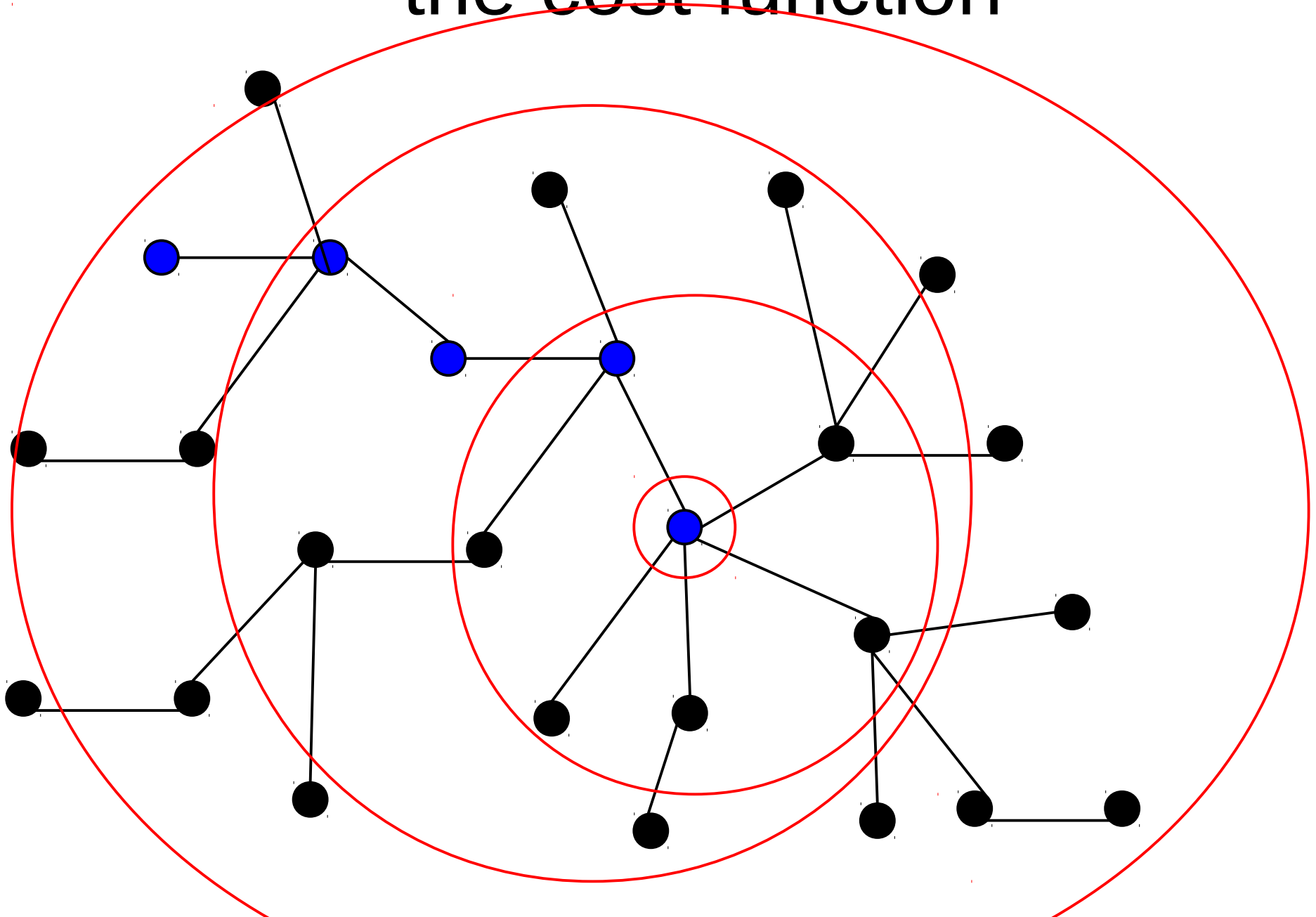
```
procedure best first  
A  $\leftarrow$   $\{s_0\}$   
while A  $\neq$   $\emptyset$   
    x  $\leftarrow$  popPriorityList(A)  
    Y  $\leftarrow$  neighbors(x)  
    A  $\leftarrow$  A  $\cup$  Y
```

Sequence of visited nodes

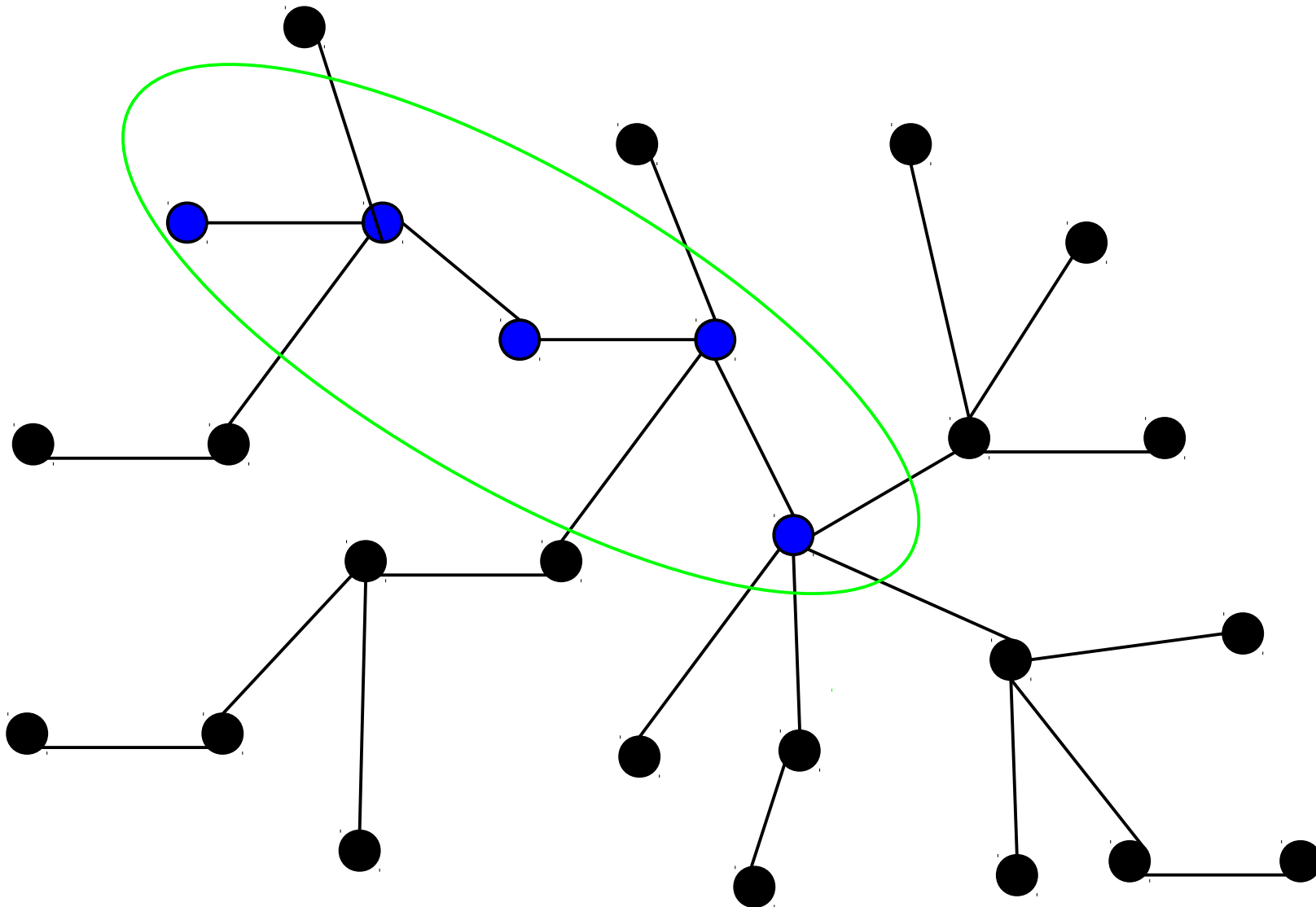




Best first based on the cost function



Best first based on an oracle

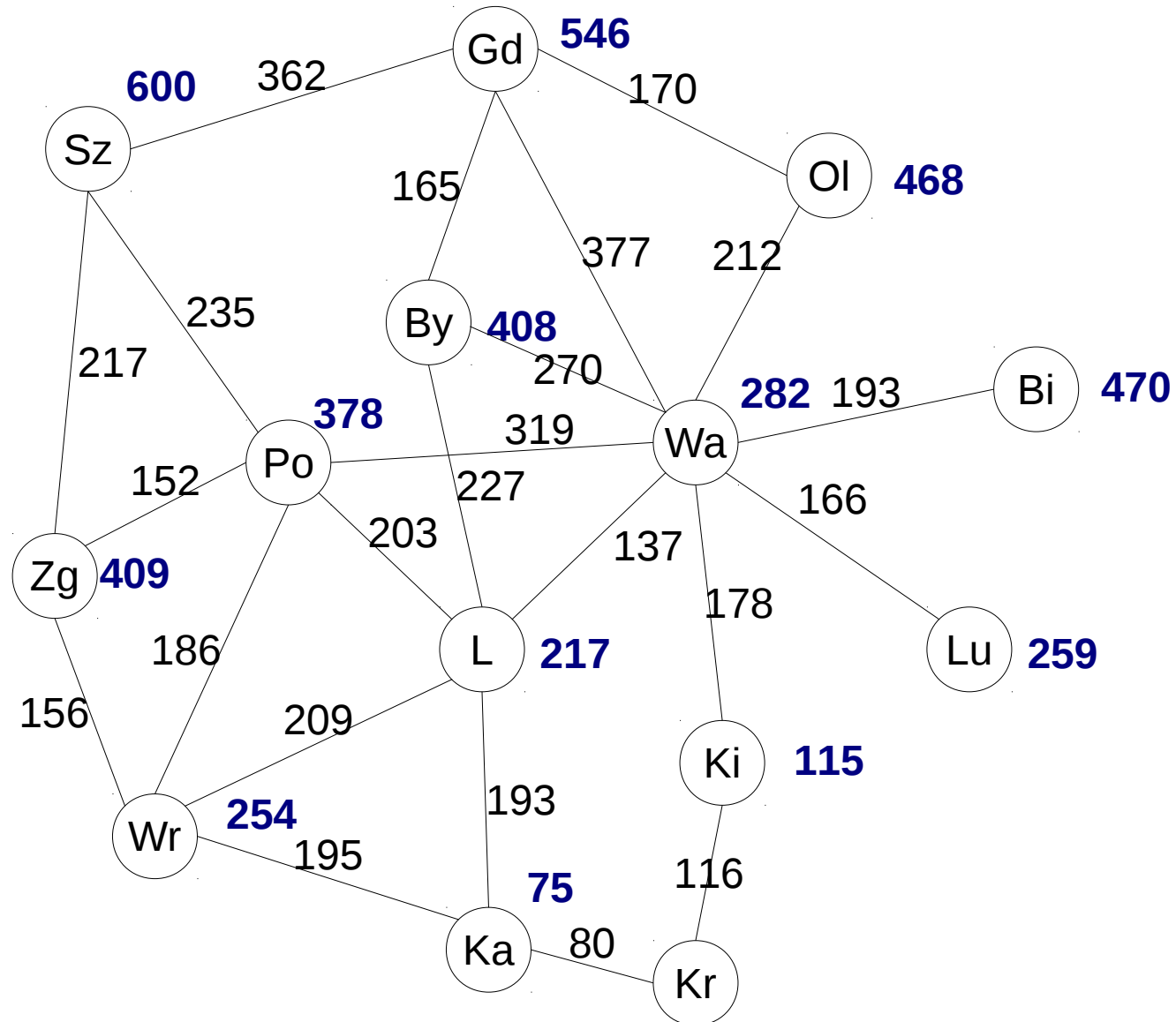


WANTED



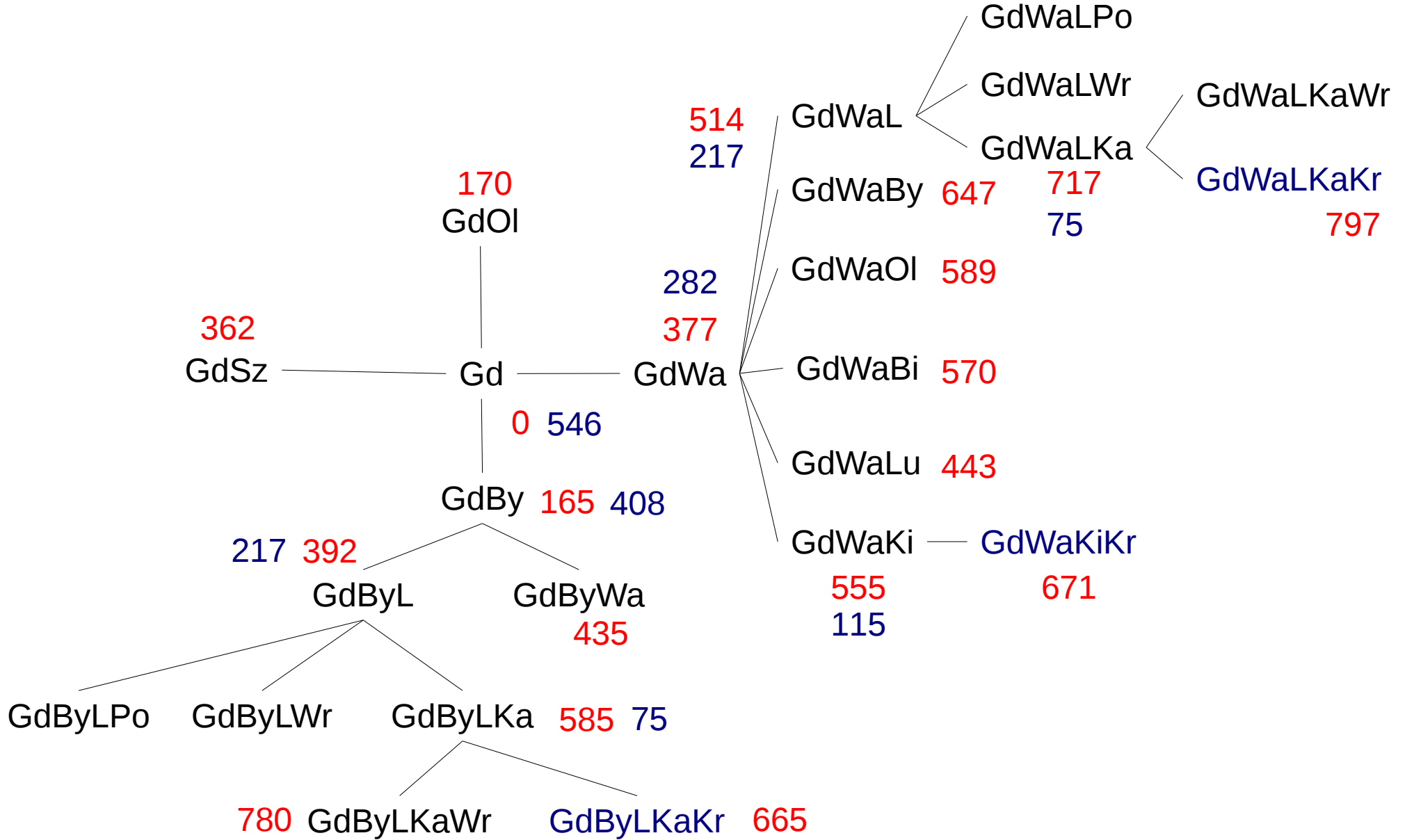
an oracle

Distances and their estimates

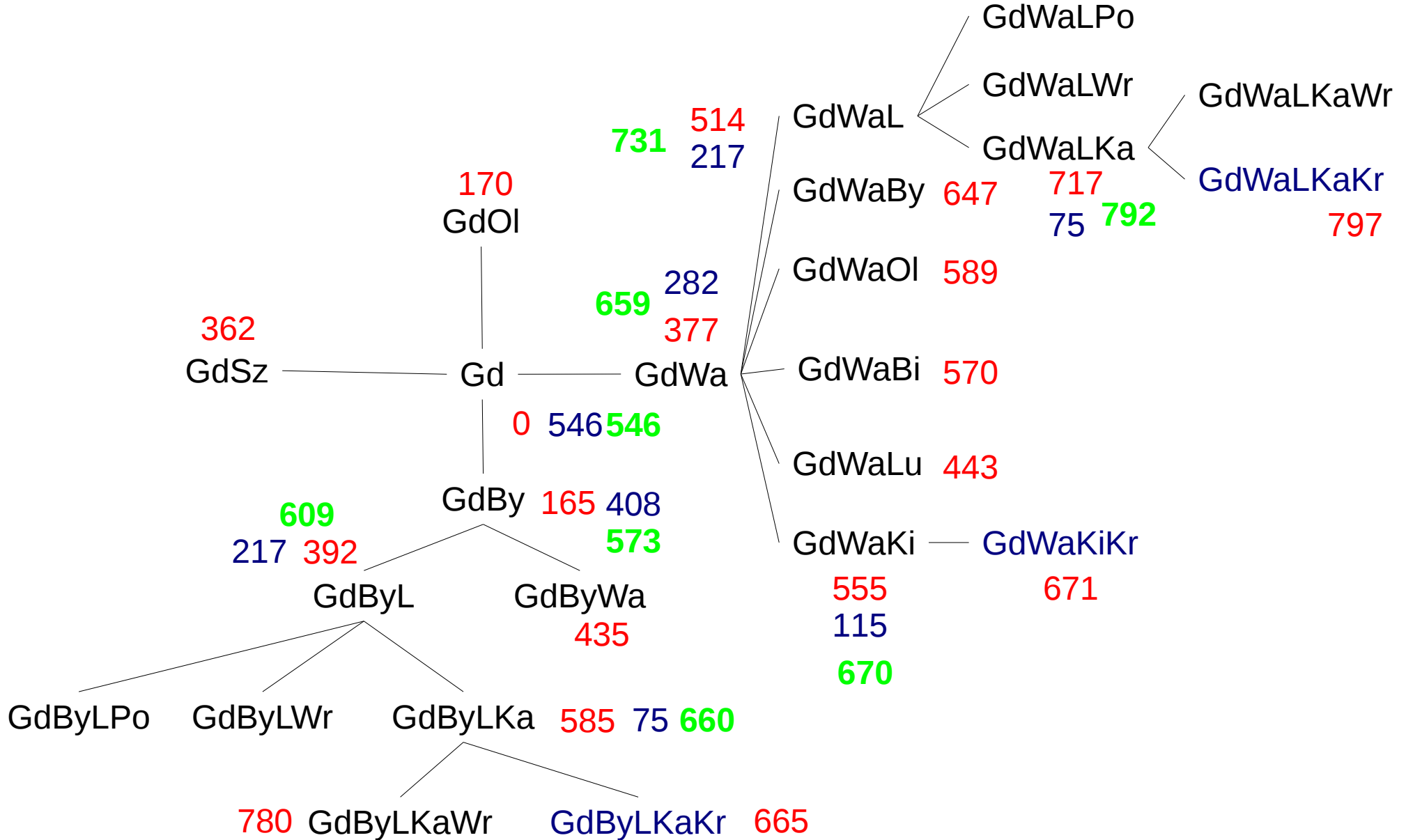


Straight line distances to Kraków

Fragment of the search space



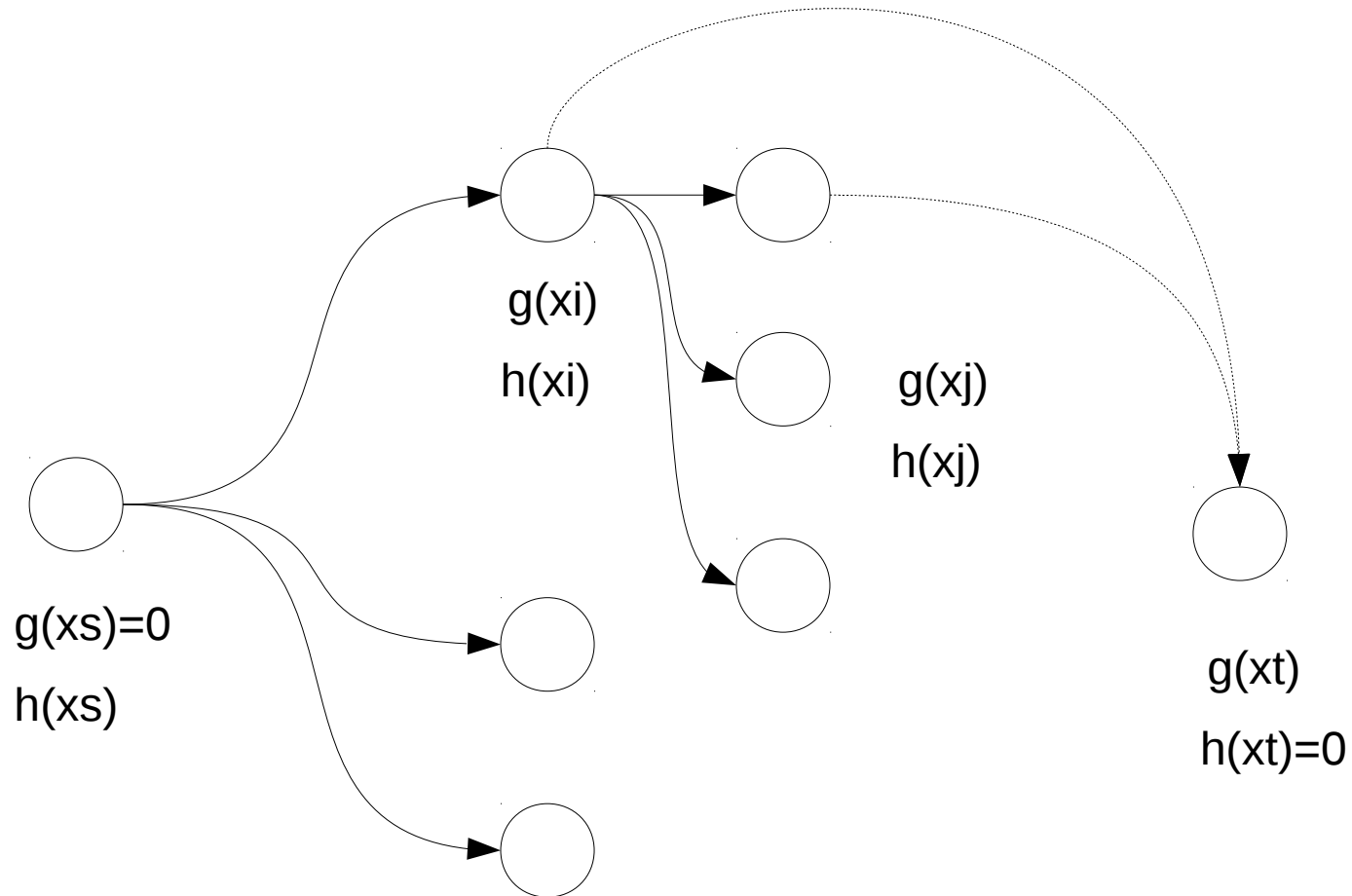
Fragment of the search space



Heuristic function

- Tree structured search space
- Objective function for intermediate nodes needed
- Cost for the intermediate node is known
- Heuristic function **estimates** how much **the cost** will **increase** after the best accessible terminal node is attained

Heuristic function (for minimization)



Optimistic (lower bound of the cost increase)

admissible: $g(x)+h(x)\leq g(x_t)$

Error of the estimate never increases while approaching to the terminal node

monotonous: $g(x_j)+h(x_j)\geq g(x_i)+h(x_i)$

Heuristic function (for minimization)

If the heuristic function is admissible and monotonic
and the solution is in a finite distance from the starting point

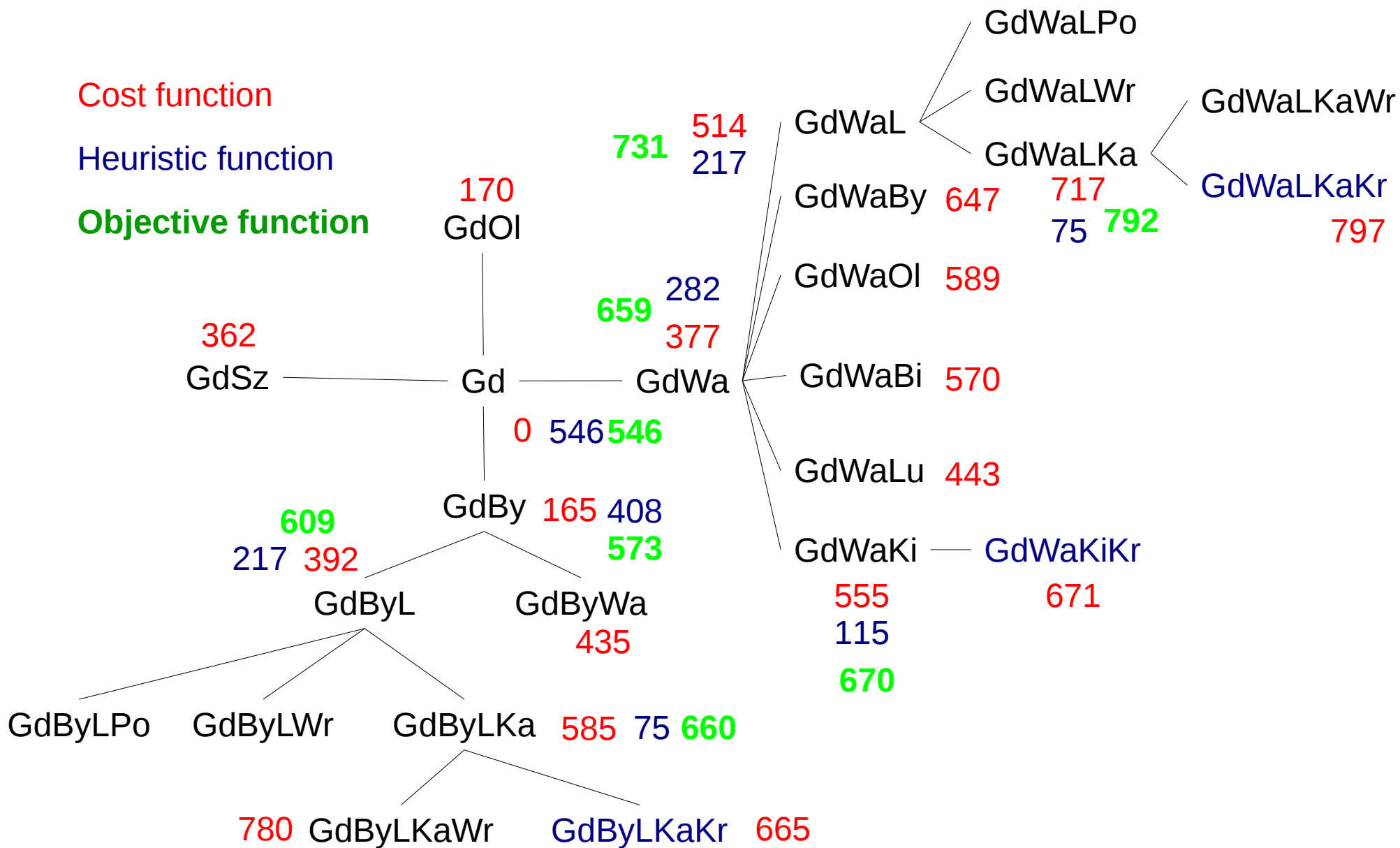
then A^* will yield that solution after a finite number of visited nodes

which is no greater than for the best-first that uses the cost function

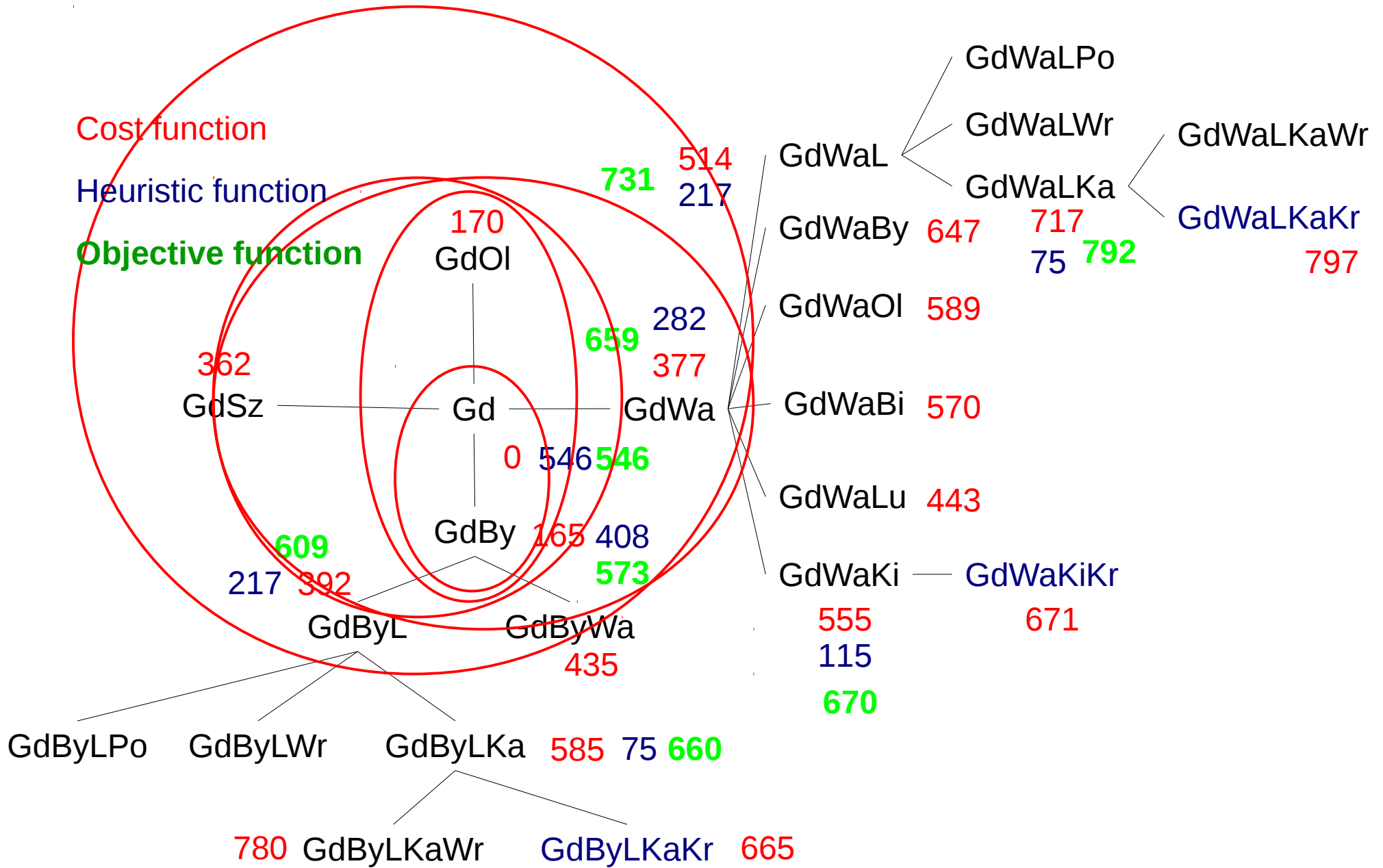
Cost function

Heuristic function

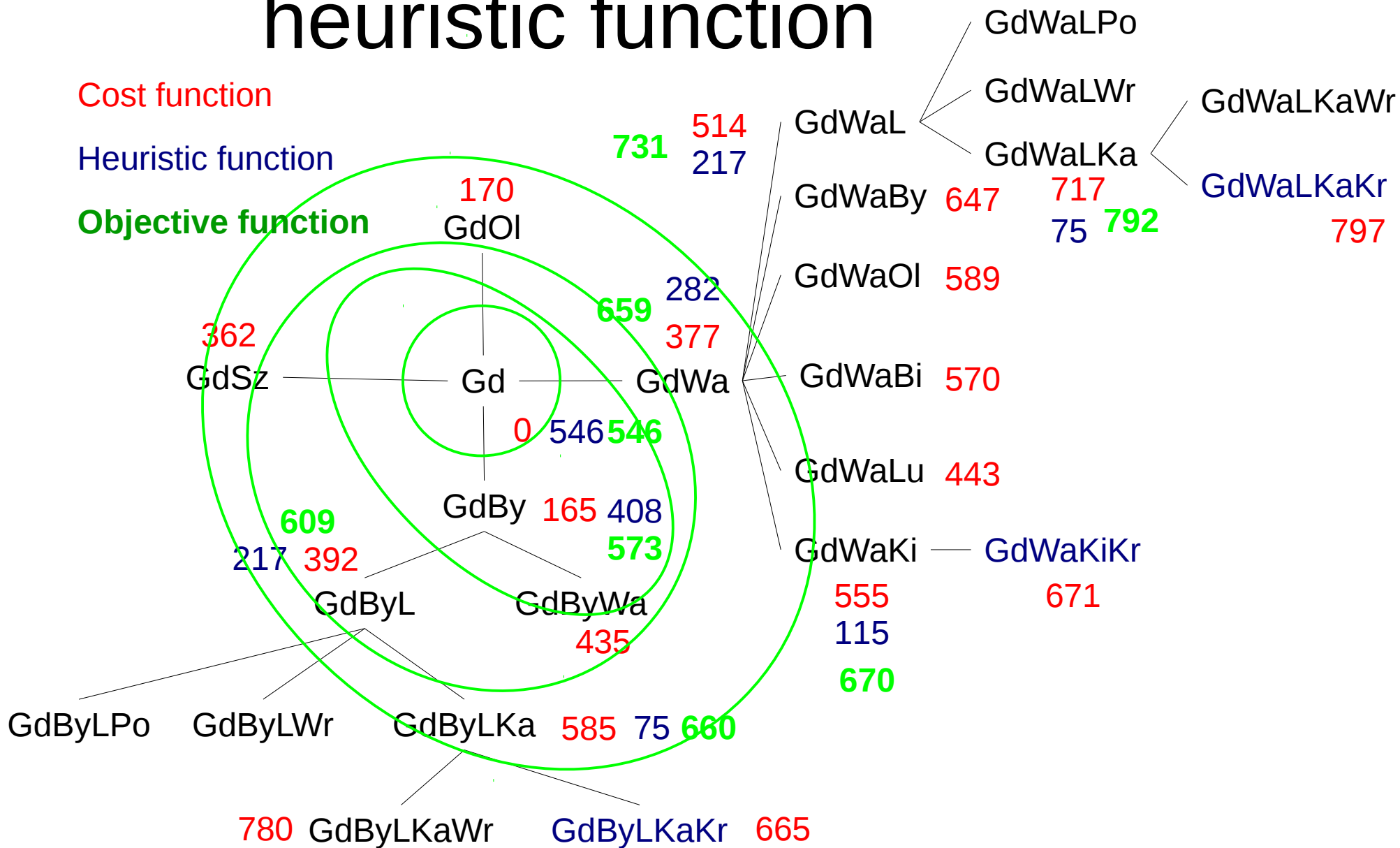
Objective function



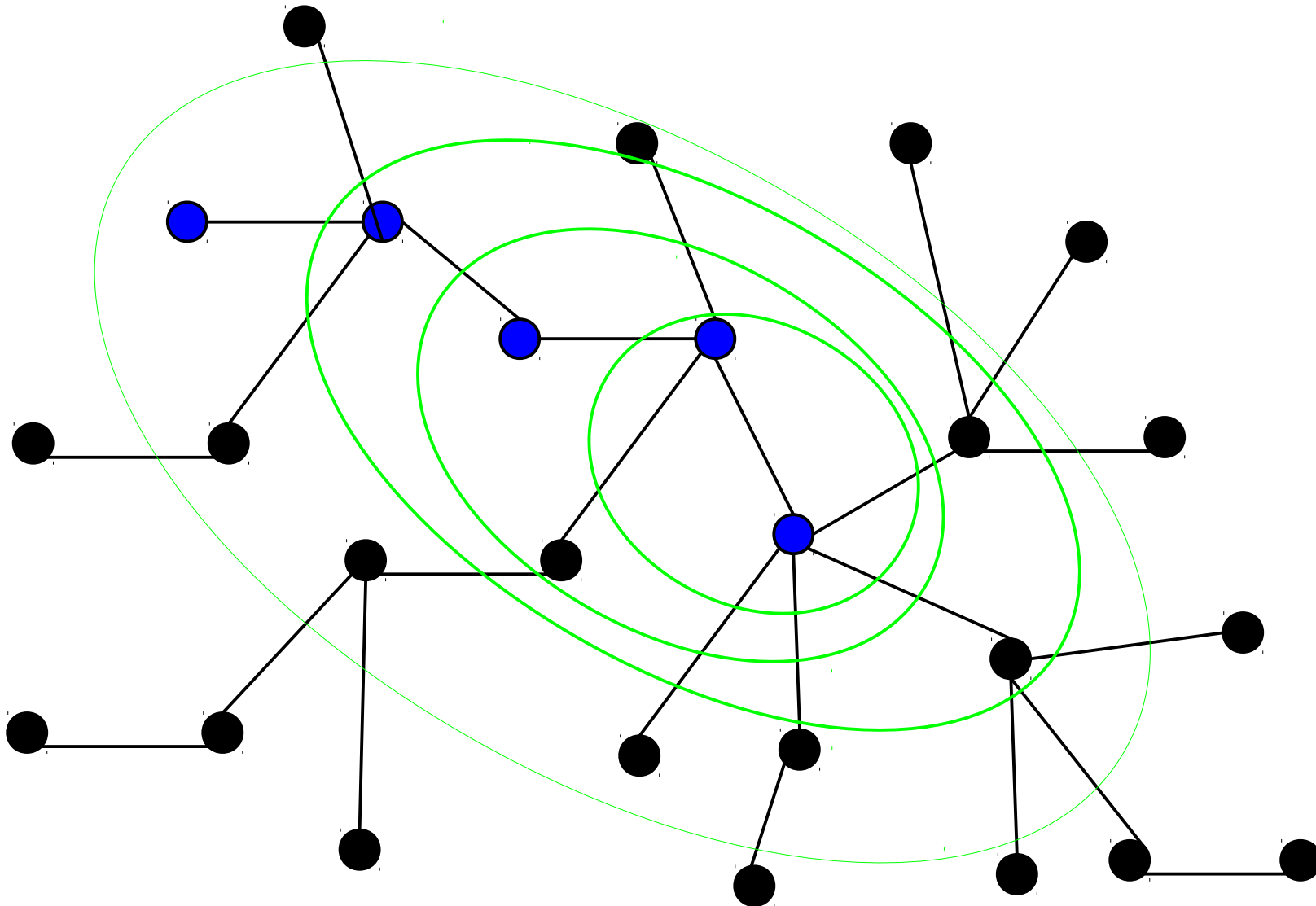
Best first based on the cost



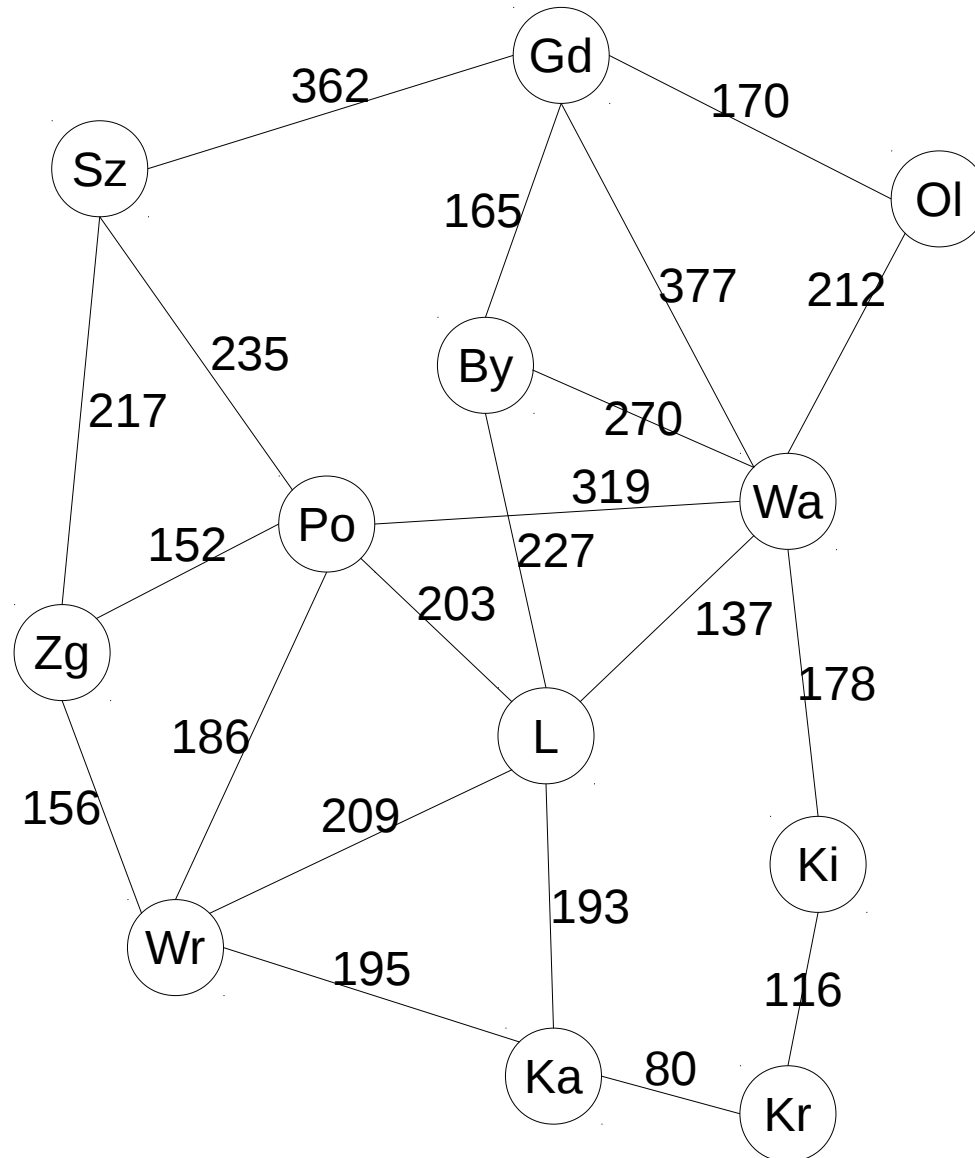
Best first based on the cost + heuristic function



Realistic heuristic function

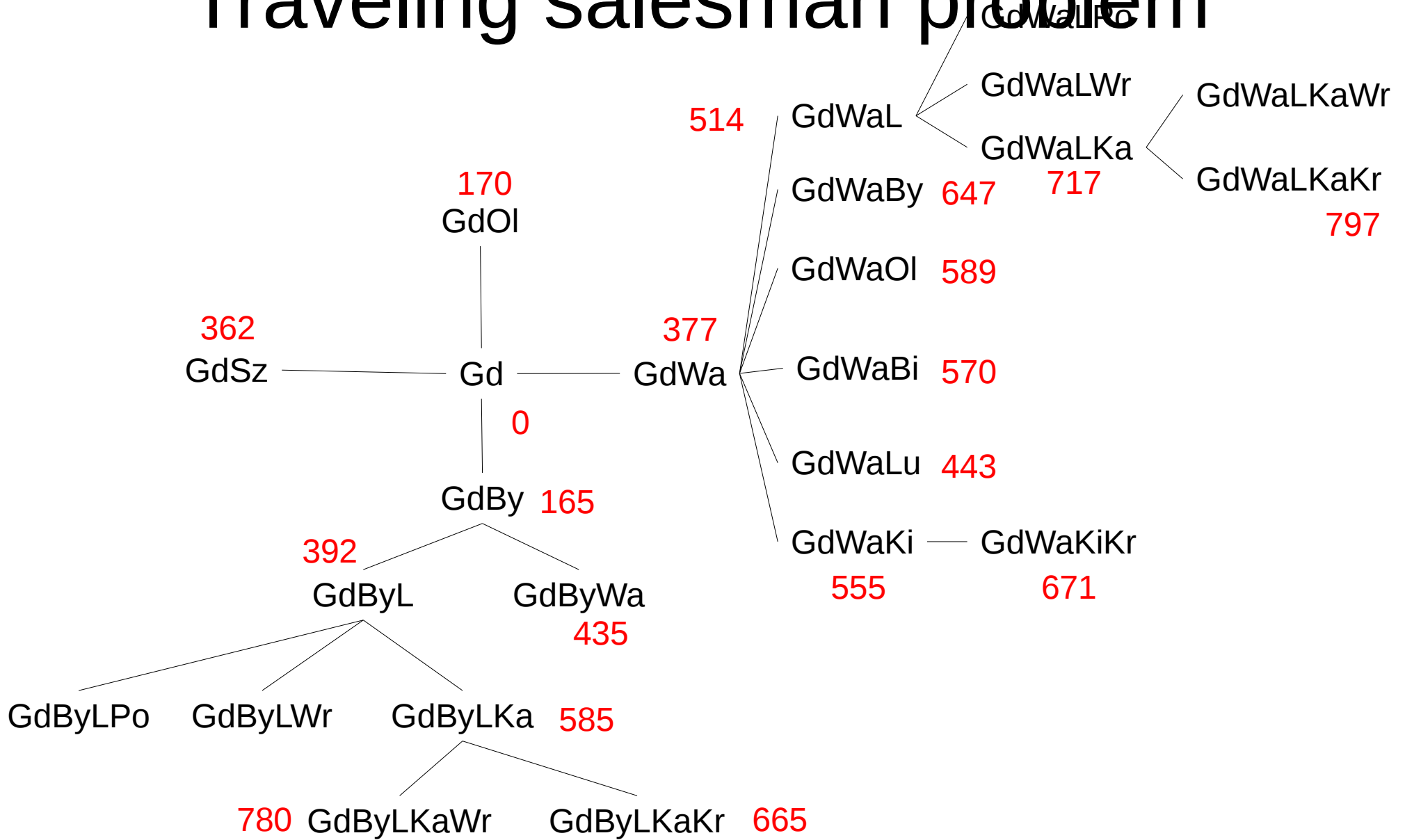


Traveling salesman problem

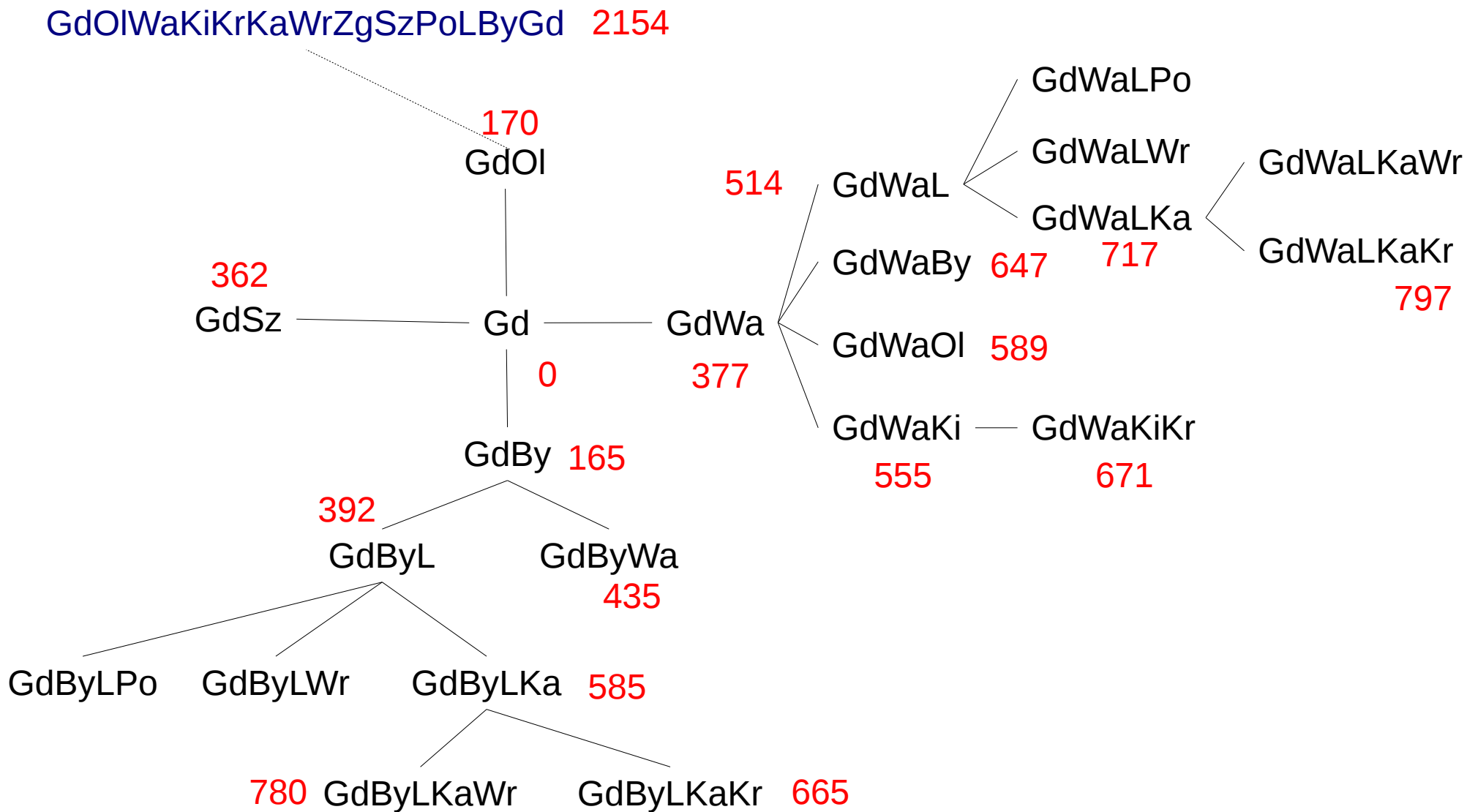


Find the
cheapest tour
to visit all cities

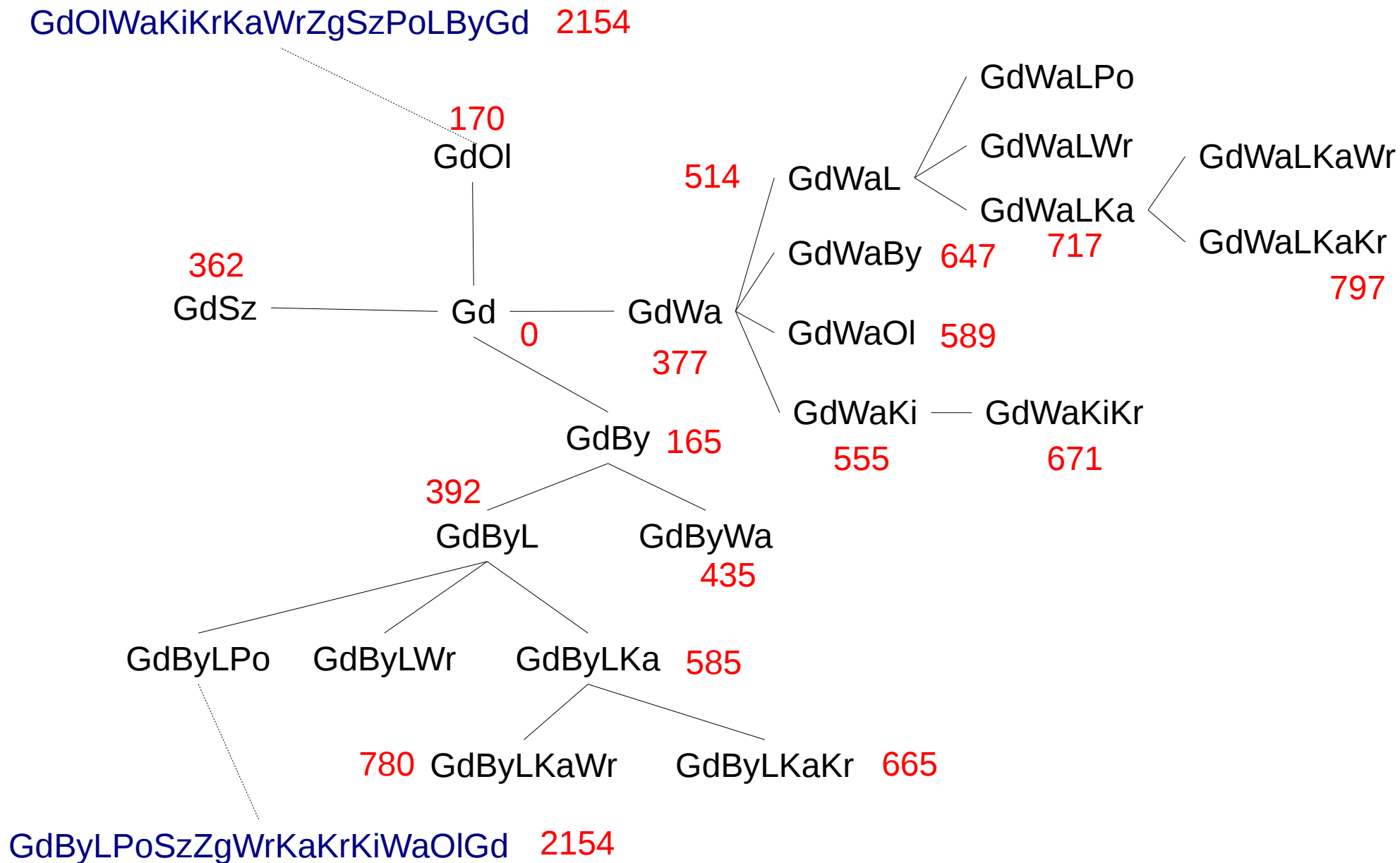
Traveling salesman problem



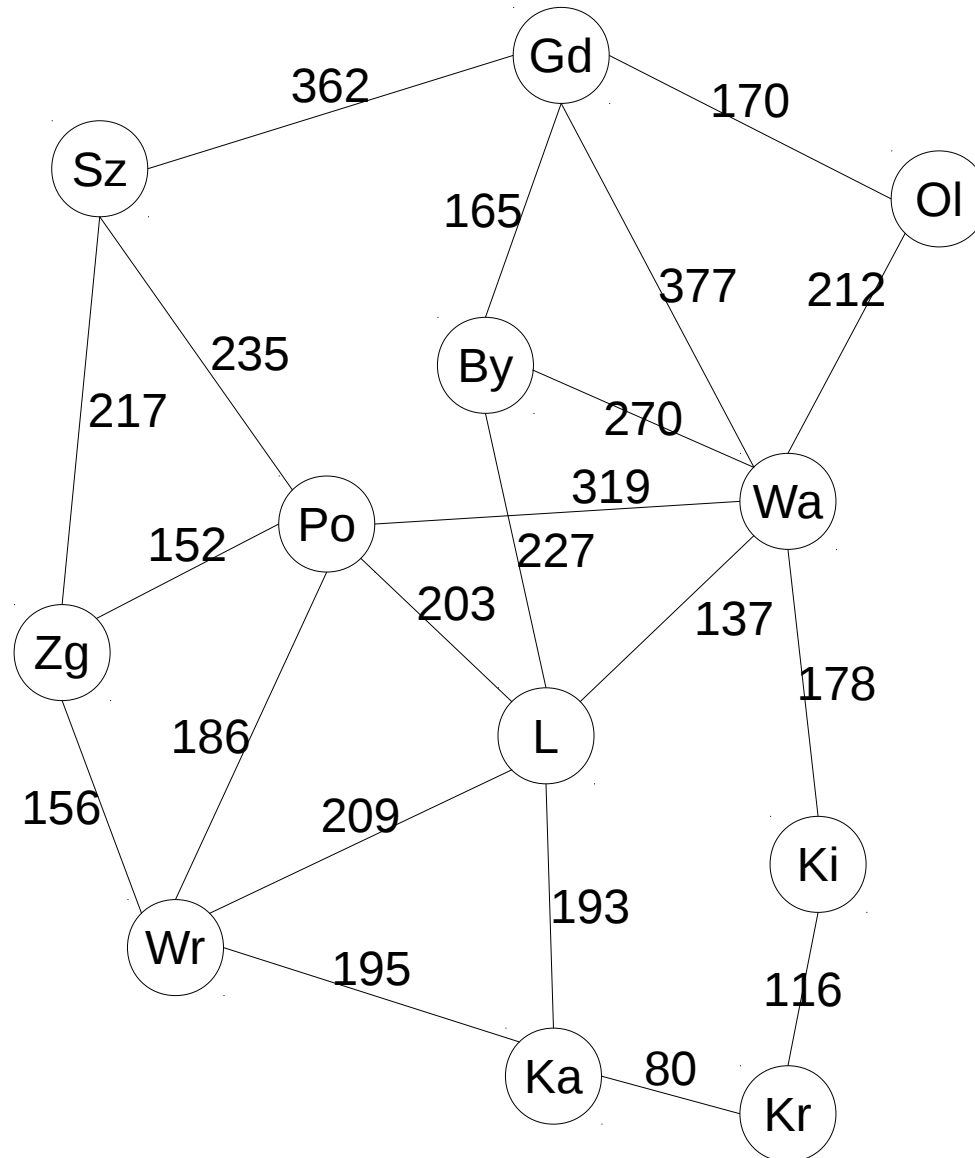
Traveling salesman problem



Traveling salesman problem

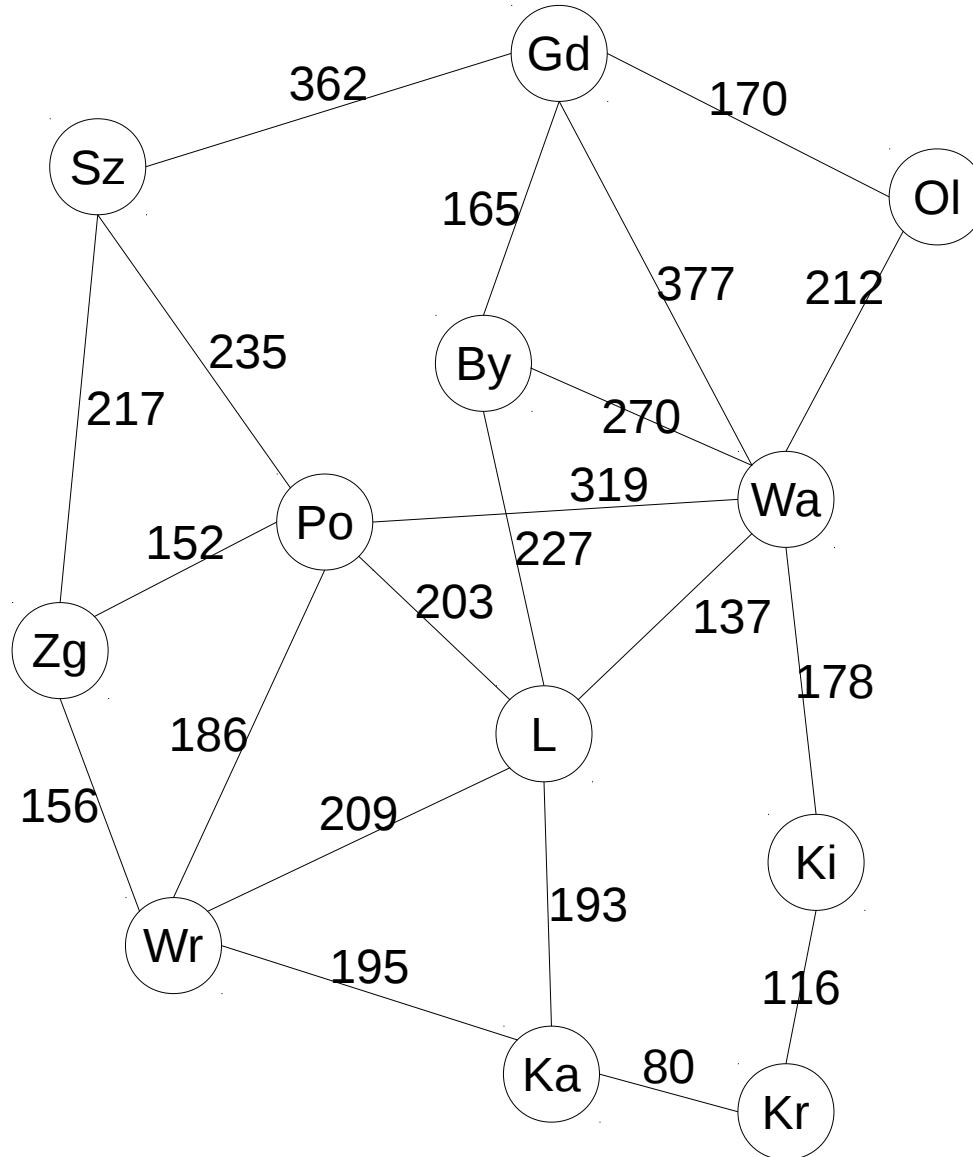


Traveling salesman problem



Heuristic
function?

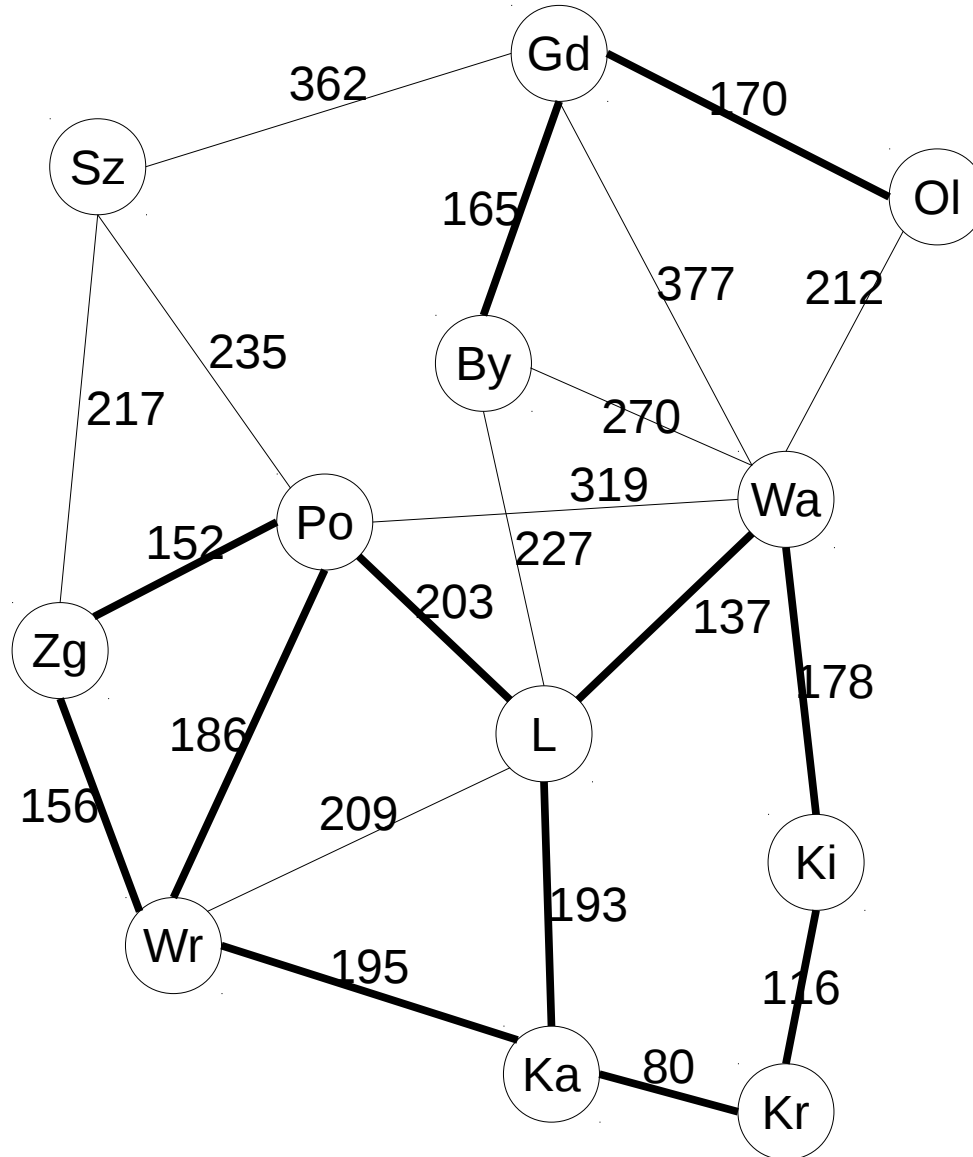
Heuristic function



Sorted list of edge weights

- 80
- 116
- 137
- 152
- 156
- 165
- 170
- 178
- 186
- 193
- 195
- 203
- 209
- 212
- 217

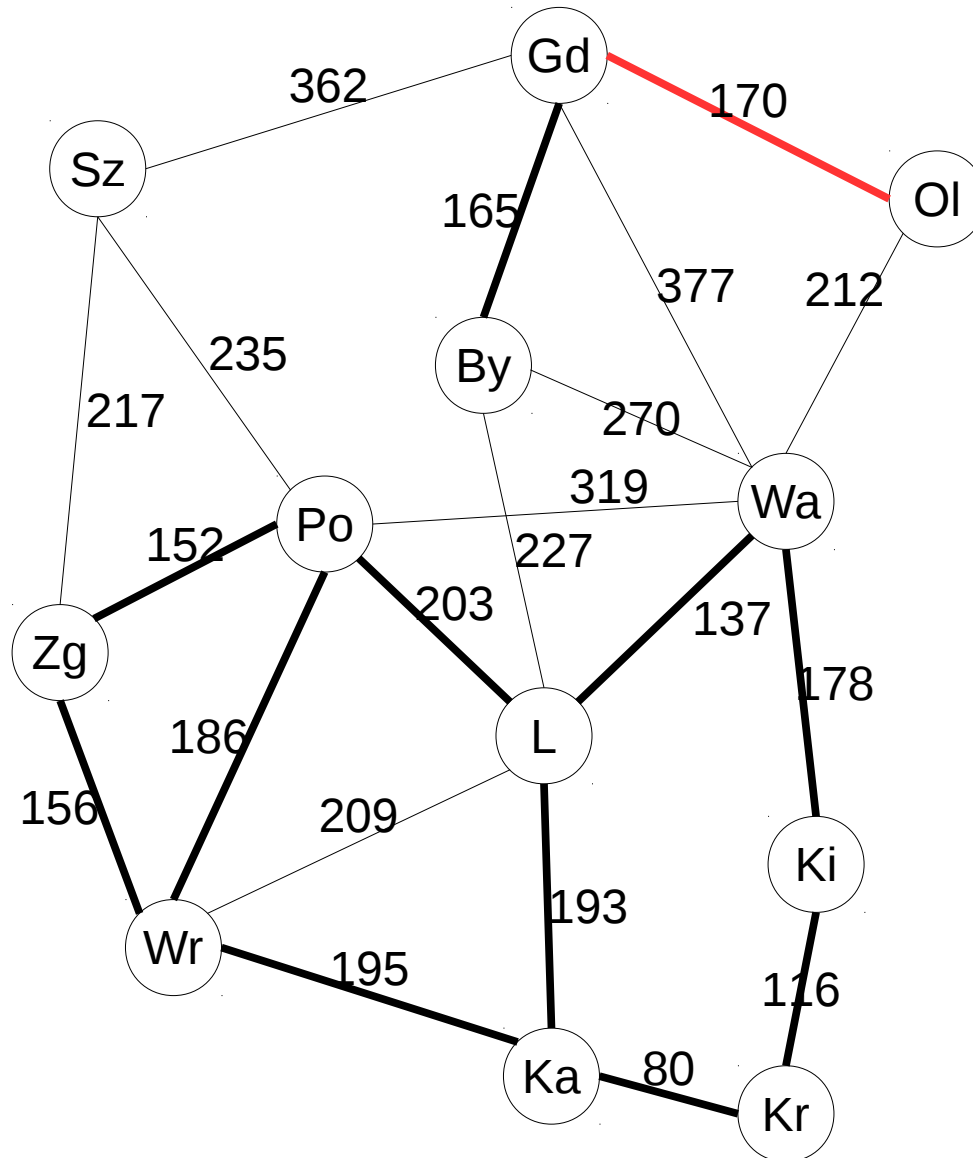
Heuristic function



Sorted list of edge weights

80
116
137
152
156
165
170
178
178
186
193
195
203
1931

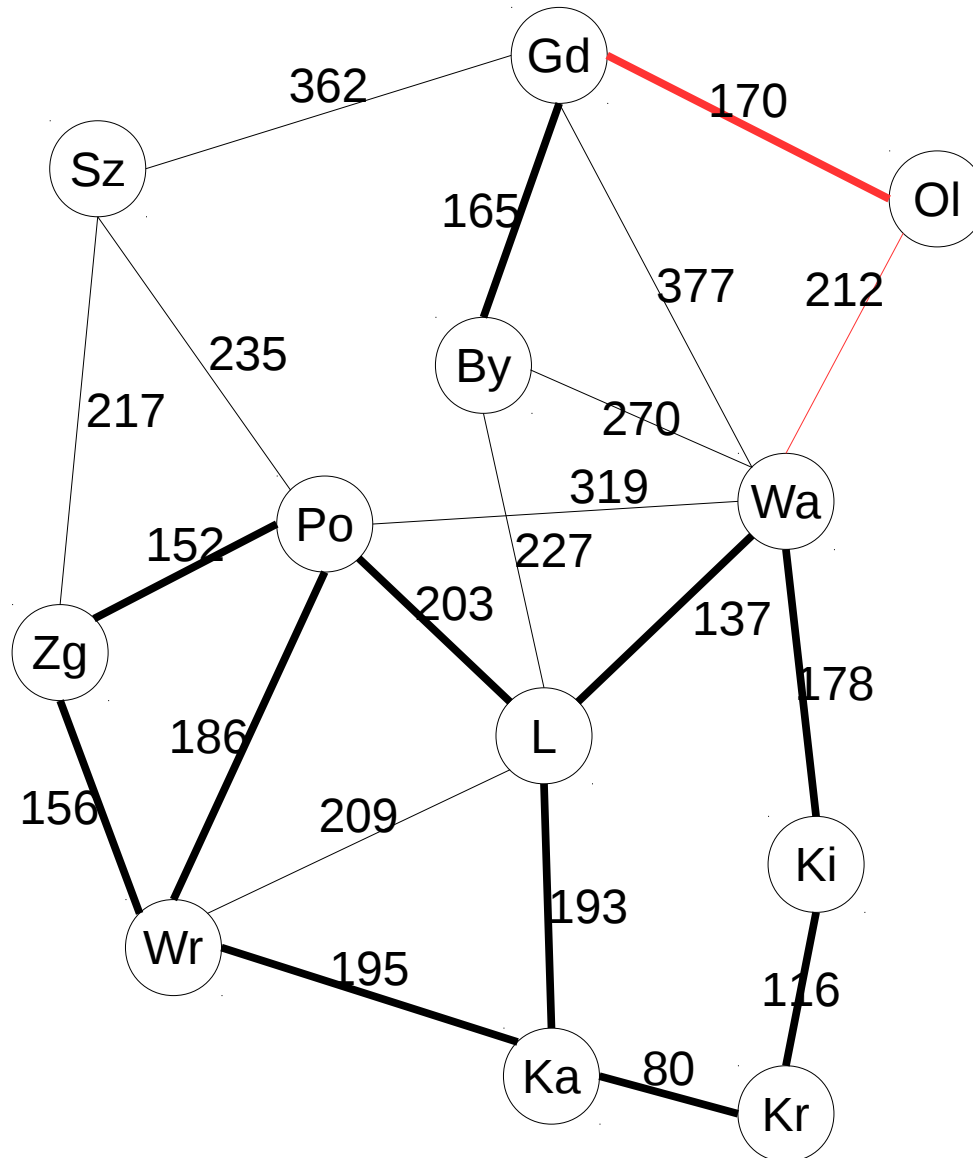
Path cost and heuristic function



Sorted list of edge weights

- 80
- 116
- 137
- 152
- 156
- 165
- 170
- 178
- 186
- 193
- 195
- 203
- 1761
- 170
- 1931

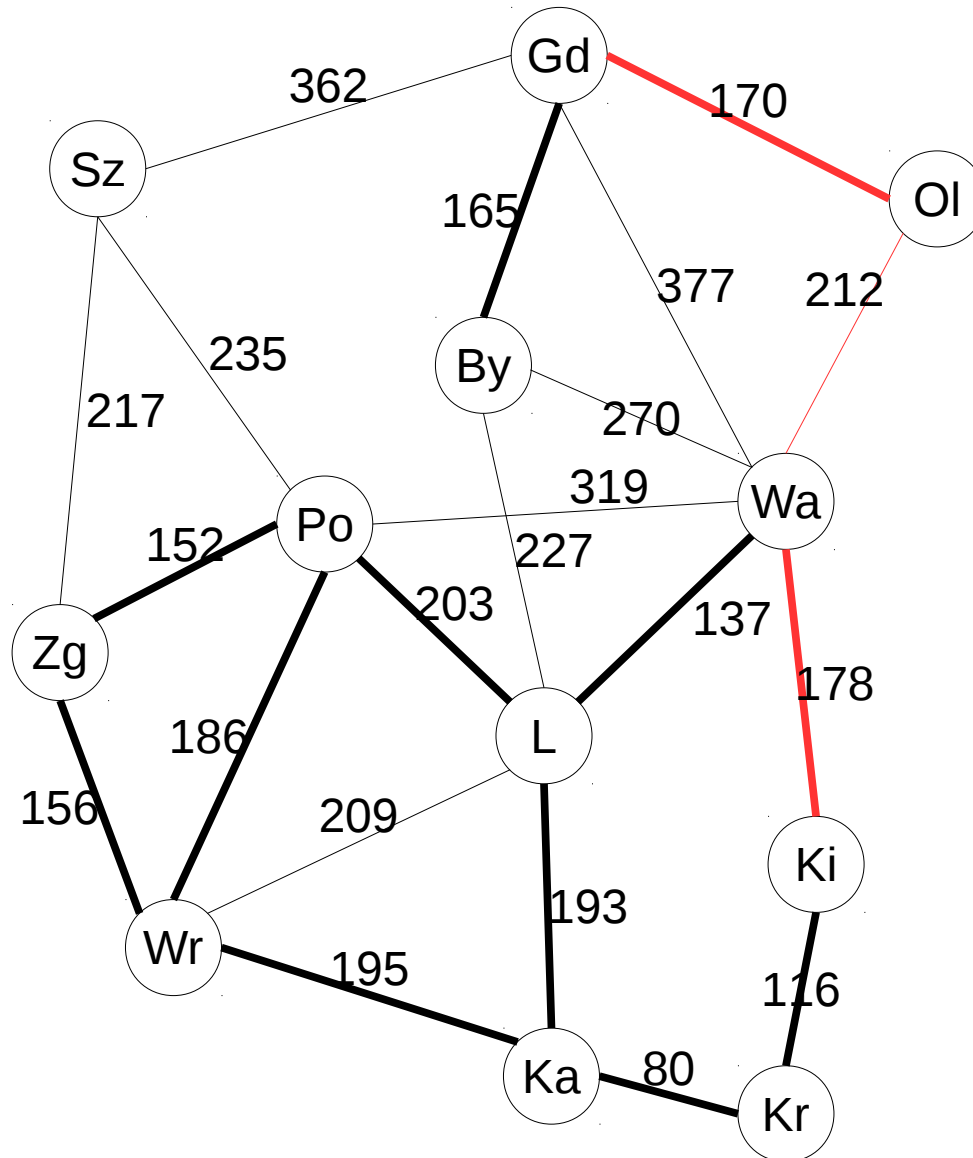
Path cost and heuristic function



Sorted list of edge weights

- 80
- 116
- 137
- 152
- 156
- 165
- 170
- 178
- 186
- 193
- 195
- 203
- 1558
- 382
- 1940

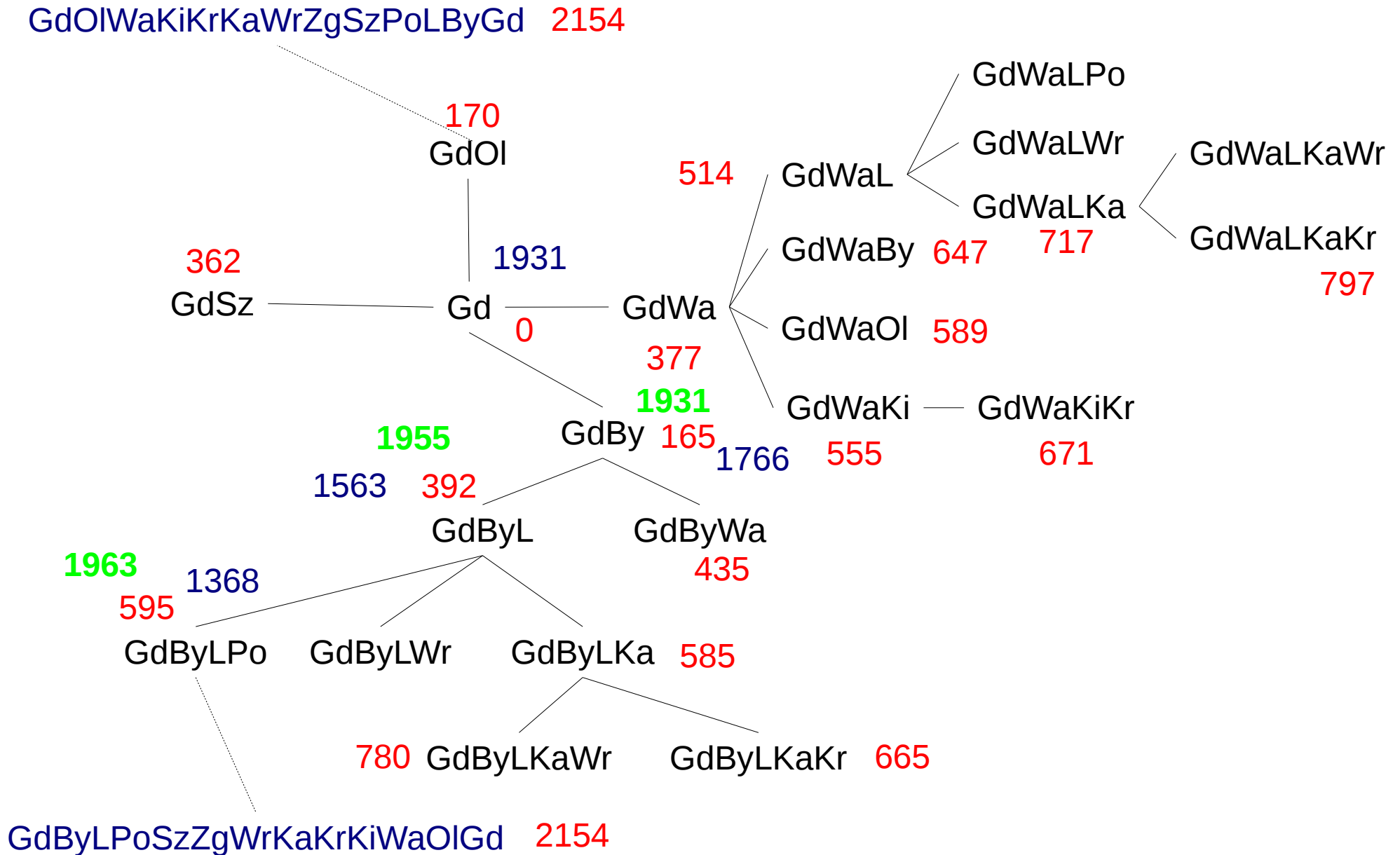
Path cost and heuristic function



Sorted list of edge weights

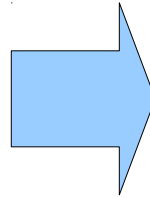
- 80
- 116
- 137
- 152
- 156
- 165
- 170
- 178
- 186
- 193
- 195
- 203
- 1380
- 560
- 1940

Traveling salesman problem



Solving 15 puzzle

1		3	12
10	11	4	13
5	6	7	14
8	9	2	15



1	2	3	4
12	13	14	5
11		15	6
10	9	8	7

Solution representation?

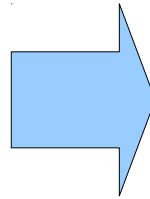
Admissible transformations?

Cost function?

Heuristic function?

Solving 15 puzzle attempt #1

1		3	12
10	11	4	13
5	6	7	14
8	9	2	15



1	2	3	4
12	13	14	5
11		15	6
10	9	8	7

Solution representation?

Positions of tiles

Admissible transformations?

Move a tile into empty place

Cost function?

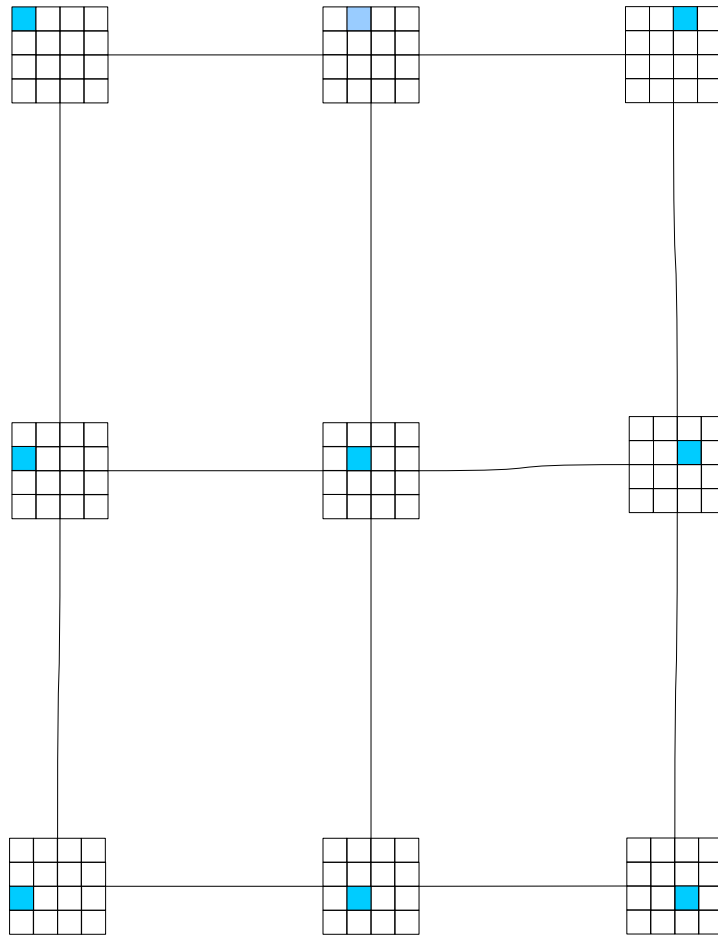
Number of wrongly placed tiles

12

Heuristic function?

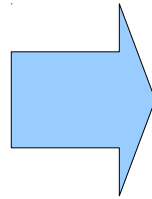
?

Solving 15 puzzle attempt #1



Solving 15 puzzle attempt #2

1		3	12
10	11	4	13
5	6	7	14
8	9	2	15



1	2	3	4
12	13	14	5
11		15	6
10	9	8	7

Solution representation?

Sequence of moves

Admissible transformations?

Move a tile into empty place

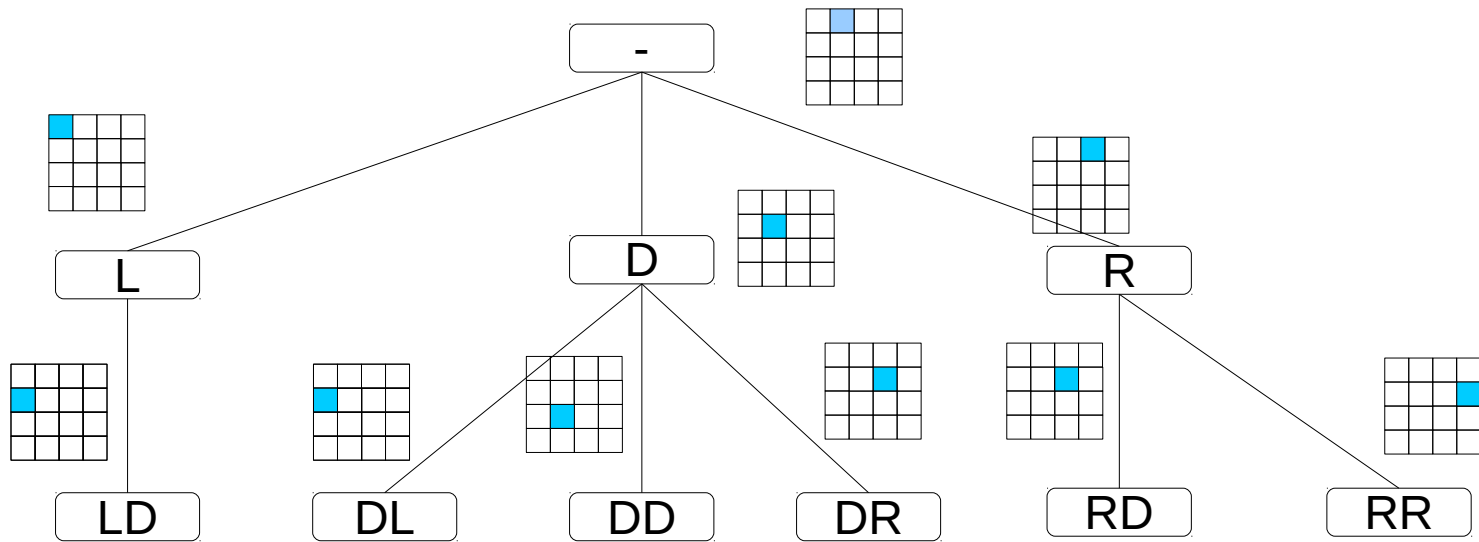
Cost function?

Number of moves

Heuristic function?

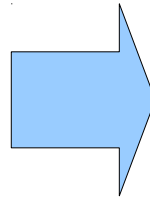
?

Solving 15 puzzle attempt #2



Solving 15 puzzle attempt #2 - heuristic function

1		3	12
10	11	4	13
5	6	7	14
8	9	2	15



1	2	3	4
12	13	14	5
11		15	6
10	9	8	7

30

2:4	10:2
4:2	11:2
5:4	12:4
6:2	13:2
7:2	14:2
8:2	15:2

Knapsack problem

- N items
- Each item has its weight $w_i > 0$ and profit $p_i > 0$
- Choose items such that total profit is maximized and total weight does not exceed W

$$\max \sum_{i=1}^n x_i p_i$$

$$\sum_{i=1}^n x_i w_i \leq W$$

$$x_i \in \{0, 1\}$$

Knapsack problem

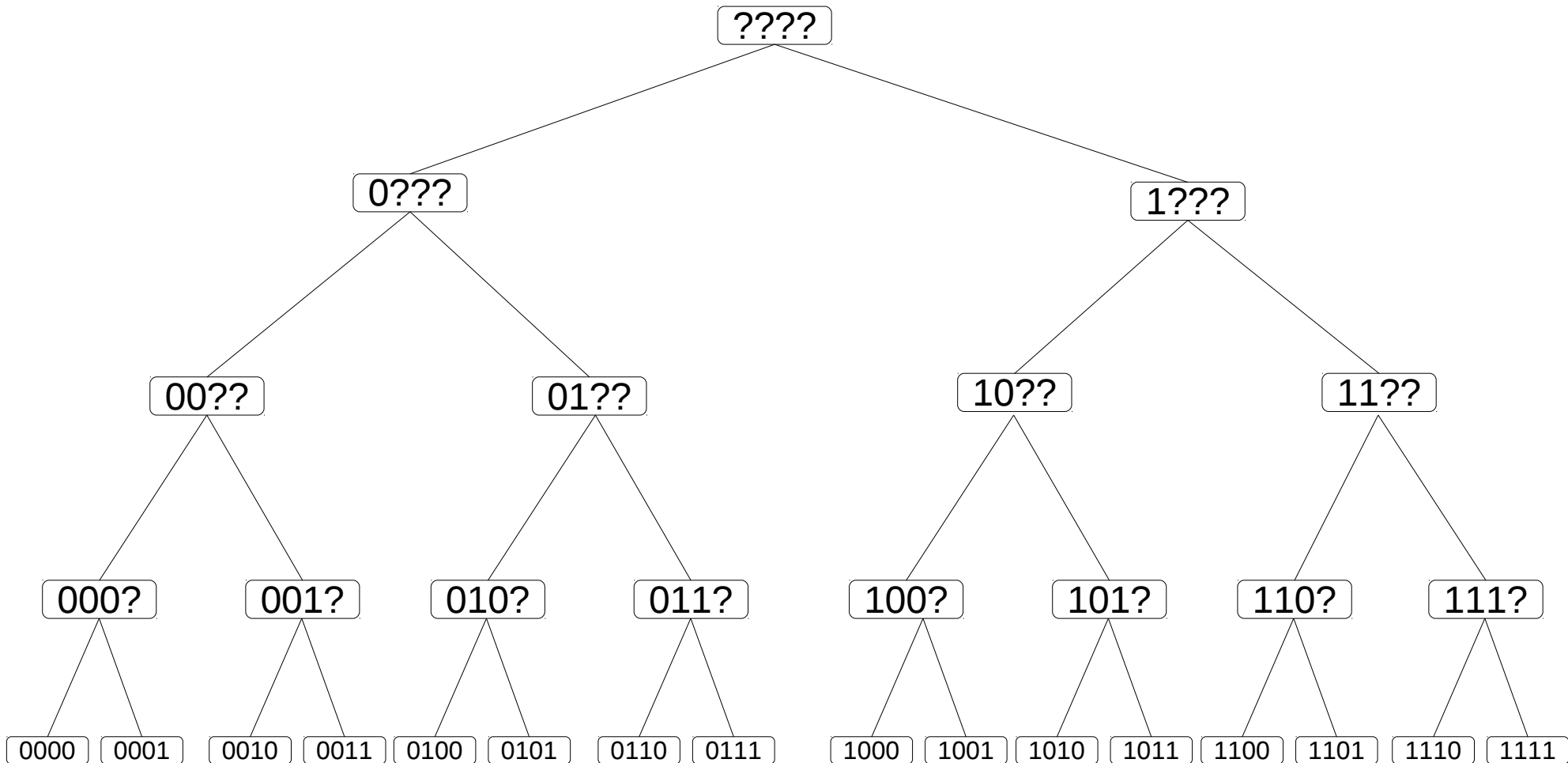
- N items
- Each item has its weight $w_i > 0$ and profit $p_i > 0$
- Choose items such that total profit is maximized and total weight does not exceed W



$$\begin{aligned} \max \quad & \sum_{i=1}^n x_i p_i \\ \sum_{i=1}^n x_i w_i & \leq W \\ x_i & \in \{0, 1\} \end{aligned}$$



Knapsack problem search space



Knapsack problem profit and heuristic function

- Profit function

$$g(x) = \sum_{i: x_i=1} p_i$$

- Heuristic function

Items are sorted w.r.t. p_i/w_i (descending)

$$h(x) = \sum_{i: x_i=?} y_i p_i$$
$$\sum_{i: x_i=?} y_i w_i = W - \sum_{i: x_i=1} x_i w_i$$
$$y_i \in [0, 1]$$

Knapsack problem example profit and heuristic function

- Items

i	1	2	3	4	5	6
p_i	20	5	10	5	6	3
w_i	6	2	5	3	4	3
p_i/w_i	3.33	2.5	2	1.66	1.5	1

- $W=13$

- Consider the solution

$$x = ? 0 ? 1 ? 1$$

- Total profit:

$$g(x) = p_4 + p_6 = 5 + 3 = 8$$

- Total weight:

$$w(x) = w_4 + w_6 = 3 + 3 = 6$$

- Vector y :

$$y = [1, 0, 1/5, 0, 0, 0]$$

- Heuristic function:

$$h(x) = 20 \cdot 1 + 10 \cdot 1/5 = 22$$